

## INTRODUCTION TO COMPUTING SYSTEMS

From Bits and Gates to C and Beyond

Yale N. Patt Sanjay J. Patel

For University of Cincinnati



#### Selected Chapters from

# INTRODUCTION TO COMPUTING SYSTEMS

### From Bits and Gates to C and Beyond

First Edition

Yale N. Patt
The University of Texas at Austin

Sanjay J. Patel University of Illinois at Urbana-Champaign

For

University of Cincinnati

The McGraw-Hill Companies, Inc. Primis Custom Publishing

New York St. Louis San Francisco Auckland Bogotá Caracas Lisbon London Madrid Mexico Milan Montreal New Delhi Paris San Juan Singapore Sydney Tokyo Toronto

## McGraw-Hill Higher Education A Division of The McGraw-Hill Companies

## Selected Chapters from Introduction to Computing Systems From Bits and Gates to C and Beyond

Copyright ©2000 by The McGraw-Hill Companies, Inc. All rights reserved. Printed in the United States of America. Except as permitted under the United States Copyright Act of 1976, no part of this publication may be reproduced or distributed in any form or by any means, or stored in a data base retrieval system, without prior written permission of the publisher.

This book contains selected material from Introduction to Computing Systesm: From Bits and Gates to C and Beyond, First edition by Yale N. Patt and Sanjay J. Patel. Copyright © 2001, by the McGraw-Hill Companies, Inc. Reprinted with permission of the publisher.

1234567890 QSR QSR 00

ISBN 0-07-244100-3

Custom Editor: Judith Wetherington

Printer/Binder: Quebecor World, Dubuque, Iowa

To the memory of my parents,
Abraham Walter Patt A''H and Sarah Clara Patt A''H,
who taught me to value "learning"
even before they taught me to ride a bicycle.

To Mira and her grandparents, Sharda Patel and Jeram Patel.

#### **PREFACE**

This textbook has evolved from EECS 100, the first computing course for computer science, computer engineering, and electrical engineering majors at the University of Michigan, that Kevin Compton and the first author introduced for the first time in the fall term, 1995.

EECS 100 happened because Computer Science and Engineering faculty had been dissatisfied for many years with the lack of student comprehension of some very basic concepts. For example, students had a lot of trouble with pointer variables. Recursion seemed to be "magic," beyond understanding.

We decided in 1993 that the conventional wisdom of starting with a high-level programming language, which was the way we (and most universities) were doing it, had its shortcomings. We decided that the reason students were not getting it was that they were forced to memorize technical details when they did not understand the basic underpinnings.

The result is the bottom-up approach taken in this book. We treat (in order) MOS transistors (very briefly, long enough for students to grasp their global switch-level behavior), logic gates, latches, logic structures (MUX, Decoder, Adder, gated latches), finally culminating in an implementation of memory. From there, we move on to the Von Neumann model of execution, then a simple computer (the LC-2), machine language programming of the LC-2, assembly language programming of the LC-2, the high level language C, recursion, pointers, arrays, and finally some elementary data structures.

We do not endorse today's popular information hiding approach when it comes to learning. Information hiding is a useful productivity enhancement technique after one understands what is going on. But until one gets to that point, we insist that information hiding gets in the way of understanding. Thus, we continually build on what has gone before, so that nothing is magic, and everything can be tied to the foundation that has already been laid.

We should point out that we do not disagree with the notion of top-down *design*. On the contrary, we believe strongly that top-down design is correct design. But there is a clear difference between how one approaches a design problem (after one understands the underlying building blocks), and what it takes to get to the point where one does understand the building blocks. In short, we believe in top-down design, but bottom-up learning for understanding.

#### WHAT IS IN THE BOOK

The book breaks down into two major segments, a) the underlying structure of a computer, as manifested in the LC-2; and b) programming in a high level language, in our case C.

#### The LC-2

We start with the underpinnings that are needed to understand the workings of a real computer. Chapter 2 introduces the bit and arithmetic and logical operations on bits. Then we begin to build the structure needed to understand the LC-2. Chapter 3 takes the student from MOS transistor, step by step, to a real memory. Our real memory consists of 4 words of 3 bits each, rather than 64 megabytes. The picture fits on a single page (Figure 3.20), making it easy for a student to grasp. By the time the students get there, they have been exposed to all the elements

that make memory work. Chapter 4 introduces the Von Neumann execution model, as a lead-in to Chapter 5, the LC-2.

The LC-2 is a 16-bit architecture that includes physical I/O via keyboard and monitor; TRAPs to the operating system for handling service calls; conditional branches on N, Z, and P condition codes; a subroutine call/return mechanism; a minimal set of operate instructions (ADD, AND, and NOT); and various addressing modes for loads and stores (direct, indirect, base+offset, and an immediate mode for loading effective addresses).

Chapter 6 is devoted to programming methodology (stepwise refinement) and debugging, and Chapter 7 is an introduction to assembly language programming. We have developed a simulator and an assembler for the LC-2. Actually, we have developed two simulators, one that runs on Windows platforms and one that runs on UNIX. The Windows simulator is available on the website and on the CD-ROM. Students who would rather use the UNIX version can download and install the software from the web at no charge.

Students use the simulator to test and debug programs written in LC-2 machine language and in LC-2 assembly language. The simulator allows on-line debugging (deposit, examine, single-step, set breakpoint, and so on). The simulator can be used for simple LC-2 machine language and assembly language programming assignments, which are essential for students to master the concepts presented throughout the first 10 chapters.

Assembly language is taught, but not to train expert assembly language programmers. Indeed, if the purpose was to train assembly language programmers, the material would be presented in an upper-level course, not in an introductory course for freshmen. Rather, the material is presented in Chapter 7 because it is consistent with the paradigm of the book. In our bottom-up approach, by the time the student reaches Chapter 7, he/she can handle the process of transforming assembly language programs to sequences of 0s and 1s. We go through the process of assembly step-by-step for a very simple LC-2 Assembler. By hand assembling, the student (at a very small additional cost in time) reinforces the important fundamental concept of translation.

It is also the case that assembly language provides a user-friendly notation to describe machine instructions, something that is particularly useful for the second half of the book. Starting in Chapter 11, when we teach the semantics of C statements, it is far easier for the reader to deal with ADD R1, R2, R3 than with 0001001010000011.

Chapter 8 deals with physical input (from a keyboard) and output (to a monitor). Chapter 9 deals with TRAPs to the operating system, and subroutine calls and returns. Students study the operating system routines (written in LC-2 code) for carrying out physical I/O invoked by the TRAP instruction.

The first half of the book concludes with Chapter 10, a treatment of stacks and data conversion at the LC-2 level, and a comprehensive example that makes use of both. The example is the simulation of a calculator, which is implemented by a main program and 11 subroutines.

#### The Language C

From there, we move on to C. The C programming language occupies the second half of the book. By the time the student gets to C, he/she has an understanding of the layers below.

The C programming language fits very nicely with our bottom-up approach. Its low-level nature allows students to see clearly the connection between software and the underlying hardware. In this book we focus on basic concepts such as control structures, functions, and arrays. Once basic programming concepts are mastered, it is a short step for students to learn more advanced concepts such as objects and abstraction.

Each time a new construct in C is introduced, the student is shown the LC-2 code that a compiler would produce. We cover the basic constructs of C (variables, operators, control, functions), pointers, recursion, arrays, structures, I/O, complex data structures, and dynamic allocation.

Chapter 11 is a gentle introduction to high-level programming languages. At this point, students have dealt heavily with assembly language and can understand the motivation behind what high-level programming languages provide. Chapter 11 also contains a simple C program, which we use to kick-start the process of learning C.

Chapter 12 deals with values, variables, constants, and operators. Chapter 13 introduces C control structures. We provide many complete program examples to give students a sample of how each of these concepts are used in practice. LC-2 code is used to demonstrate how each C construct affects the machine at the lower levels.

In Chapter 14, students are exposed to techniques for debugging high-level source code. Chapter 15 introduces functions in C. Students are not merely exposed to the syntax of functions. Rather they learn how functions are actually executed using a run-time stack. A number of examples are provided.

Chapter 16 teaches recursion, using the student's newly gained knowledge of functions, activation records, and the run-time stack. Chapter 17 teaches pointers and arrays, relying heavily on the students' understanding of how memory is organized. Chapter 18 introduces the details of I/O functions in C, in particular, streams, variable length argument lists, and how C I/O is affected by the various format specifications. This chapter relies on the student's earlier exposure to physical I/O in Chapter 8. Chapter 19 concludes the coverage of C with structures, dynamic memory allocation, and linked lists.

Along the way, we have tried to emphasize good programming style and coding methodology by means of examples. Novice programmers probably learn at least as much from the programming examples they read as from the rules they are forced to study. Insights that accompany these examples are highlighted by means of lightbulb icons that are included in the margins.

We have found that the concept of pointer variables (Chapter 17) is not at all a problem. By the time students encounter it, they have a good understanding of what memory is all about, since they have analyzed the logic design of a small memory (Chapter 3). They know the difference, for example, between a memory location's address and the data stored there.

Recursion ceases to be magic since, by the time a student gets to that point (Chapter 16), he/she has already encountered all the underpinnings. Students understand how stacks work at the machine level (Chapter 10), and they understand the call/return mechanism from their LC-2 machine language programming experience, and the need for linkages between a called program and the return to the caller (Chapter 9). From this foundation, it is not a large step to explain functions by introducing run-time activation records (Chapter 15), with a lot of the mystery about argument passing, dynamic declarations, and so on, going away. Since a function can call a function, it is one additional small step (certainly no magic involved) for a function to call itself.

#### How to Use this Book

We have discovered over the past two years that there are many ways the material in this book can be presented in class effectively. We suggest six presentations below.

The Michigan model. First course, no formal prerequisites. Very intensive, this course
covers the entire book. We have found that with talented, very highly motivated students,
this works best.



- Normal usage. First course, no prerequisites. This course is also intensive, although less so. It covers most of the book, leaving out Sections 10.3 and 10.4 of Chapter 10, Chapters 16 (recursion), 18 (the details of C I/O), and 19 (data structures).
- 3. Second course. Several schools have successfully used the book in their second course, after the students have been exposed to programming with an object-oriented programming language in a milder first course. In this second course, the entire book is covered, spending the first two-thirds of the semester on the first 10 chapters, and the last one-third of the semester on the second half of the book. The second half of the book can move more quickly, given that it follows both Chapter 1–10 and the introductory programming course, which the student has already taken. Since students have experience with programming, lengthier programming projects can be assigned. This model allows students who were introduced to programming via an object-oriented language to pick up C, which they will certainly need if they plan to go on to advanced software courses such as operating systems.
- Two quarters. An excellent use of the book. No prerequisites, the entire book can be covered easily in two quarters, the first quarter for Chapters 1-10, the second quarter for Chapters 11-19.
- 5. Two semesters. Perhaps the optimal use of the book. A two-semester sequence for freshmen. No formal prerequisites. First semester, Chapters 1–10, with supplemental material from Appendix C, the Microarchitecture of the LC-2. Second semester, Chapters 11–19 with additional substantial programming projects so that the students can solidify the concepts they learn in lectures.
- 6. A sophomore course in computer hardware. Some universities have found the book useful for a sophomore level breadth-first survey of computer hardware. They wish to introduce students in one semester to number systems, digital logic, computer organization, machine language and assembly language programming, finishing up with the material on stacks, activation records, recursion, and linked lists. The idea is to tie the hardware knowledge the students have acquired in the first part of the course to some of the harder to understand concepts that they struggled with in their freshman programming course. We strongly believe the better paradigm is to study the material in this book before tackling an object-oriented language. Nonetheless, we have seen this approach used successfully, where the sophomore student gets to understand the concepts in this course, after struggling with them during the freshman year.

#### Some Observations

Understanding, not Memorizing Since the course builds from the bottom up, we have found that less memorization of seemingly arbitrary rules is required than in traditional programming courses. Students understand that the rules make sense since by the time a topic is taught, they have an awareness of how that topic is implemented at the levels below it. This approach is good preparation for later courses in design, where understanding of and insights gained from fundamental underpinnings is essential to making the required design tradeoffs.

The Student Debugs the Student's Program We hear complaints from industry all the time about CS graduates not being able to program. Part of the problem is the helpful teaching assistant, who contributes far too much of the intellectual component of the student's program, so the student never has to really master the art. Our approach is to push the student to do the job without the teaching assistant (TA). Part of this comes from the bottom-up approach where memorizing is minimized and the student builds on what he/she already knows. Part of this is

the simulator, which the student uses from day one. The student is taught debugging from the beginning and is required to use the debugging tools of the simulator to get his/her programs to work from the very beginning. The combination of the simulator and the order in which the subject material is taught results in students actually debugging their own programs instead of taking their programs to the TA for help . . . and the common result that the TAs end up writing the programs for the students.

**Preparation for the Future: Cutting Through Protective Layers** In today's real world, professionals who use computers in systems but remain ignorant of what is going on underneath are likely to discover the hard way that the effectiveness of their solutions is impacted adversely by things other than the actual programs they write. This is true for the sophisticated computer programmer as well as the sophisticated engineer.

Serious programmers will write more efficient code if they understand what is going on beyond the statements in their high-level language. Engineers, and not just computer engineers, are having to interact with their computer systems today more and more at the device or pin level. In systems where the computer is being used to sample data from some metering device such as a weather meter or feedback control system, the engineer needs to know more than just how to program in FORTRAN. This is true of mechanical, chemical, and aeronautical engineers today, not just electrical engineers. Consequently, the high-level programming language course, where the compiler protects the student from everything "ugly" underneath, does not serve most engineering students well, and certainly does not prepare them for the future.

Rippling Effects Through the Curriculum The material of this text clearly has a rippling effect on what can be taught in subsequent courses. Subsequent programming courses can not only assume the students know the syntax of C but also understand how it relates to the underlying architecture. Consequently, the focus can be on problem solving and more sophisticated data structures. On the hardware side, a similar effect is seen in courses in digital logic design and in computer organization. Students start the logic design course with an appreciation of what the logic circuits they master are good for. In the computer organization course, the starting point is much further along than when students are seeing the term Program Counter for the first time. Feedback from Michigan faculty members in the follow-on courses have noticed substantial improvement in student's comprehension, compared to what they saw before students took EECS 100.

#### ACKNOWLEDGMENTS

This book has benefited greatly from important contributions of many, many people. At the risk of leaving out some, we would at least like to acknowledge the following.

First, Professor Kevin Compton. Kevin believed in the concept of the book since it was first introduced at a curriculum committee meeting that he chaired at Michigan in 1993. The book grew out of a course (EECS 100) that he and the first author developed together, and co-taught the first three semesters it was offered at Michigan in fall 1995, winter 1996, and fall 1996. Kevin's insights into programming methodology (independent of the syntax of the particular language) provided a sound foundation for the beginning student. The course at Michigan and this book would be a lot less were it not for Kevin's influence.

Several other students and faculty at Michigan were involved in the early years of EECS 100 and the early stages of the book. We are particularly grateful for the help of Professor David Kieras, Brian Hartman, David Armstrong, Matt Postiff, Dan Friendly, Rob Chappell, David Cybulski, Sangwook Kim, Don Winsor, and Ann Ford.

We also benefited enormously from TAs who were committed to helping students learn. The focus was always on how to explain the concept so the student gets it. We acknowledge, in particular, Fadi Aloul, David Armstrong, David Baker, Rob Chappell, David Cybulski, Amolika Gurujee, Brian Hartman, Sangwook Kim, Steve Maciejewski, Paul Racunas, David Telehowski, Francis Tseng, Aaron Wagner, and Paul Watkins.

We were delighted with the response from the publishing world to our manuscript. We ultimately decided on McGraw-Hill in large part because of the editor, Betsy Jones. Once she checked us out, she became a strong believer in what we are trying to accomplish. Throughout the process, her commitment and energy level have been greatly appreciated. We also appreciate what Michelle Flomenhoft has brought to the project. It has been a pleasure to work with her.

Our book has benefited from extensive reviews provided by faculty members at many universities. We gratefully acknowledge reviews provided by Carl D. Crane III, Florida, Nat Davis, Virginia Tech, Renee Elio, University of Alberta, Kelly Flangan, BYU, George Friedman, UIUC, Franco Fummi, Universita di Verona, Dale Grit, Colorado State, Thor Gulsrud, Stavanger College, Brad Hutchings, BYU, Dave Kaeli, Northeastern, Rasool Kenarangui, UT at Arlington, Joel Kraft, Case Western Reserve, Wei-Ming Lin, UT at San Antonio, Roderick Loss, Montgomery College, Ron Meleshko, Grant MacEwan Community College, Andreas Moshovos, Northwestern, Tom Murphy, The Citadel, Murali Narayanan, Kansas State, Carla Purdy, Cincinnati, T. N. Rajashekhara, Camden County College, Nello Scarabottolo, Universita degli Studi di Milano, Robert Schaefer, Daniel Webster College, Tage Stabell-Kuloe, University of Tromsoe, Jean-Pierre Steger, Burgdorf School of Engineering, Bill Sverdlik, Eastern Michigan, John Tronto, St. Michael's College, Murali Varansi, University of South Florida, Montanez Wade, Tennessee State, and Carl Wick, US Naval Academy.

In addition to all these people, there were others who contributed in many different and sometimes unique ways. Space dictates that we simply list them and say thank you. Susan Kornfield, Ed DeFranco, Evan Gsell, Rich Belgard, Tom Conte, Dave Nagle, Bruce Shriver, Bill Sayle, Steve Lumetta, Dharma Agarwal, and David Lilia.

Finally, if you will indulge the first author a bit: This book is about developing a strong foundation in the fundamentals with the fervent belief that once that is accomplished, students can go as far as their talent and energy can take them. This objective was instilled in me by the professor who taught me how to be a professor, Professor William K. Linvill. It has been more than 35 years since I was in his classroom, but I still treasure the example he set.

#### A FINAL WORD

We hope you will enjoy the approach taken in this book. Nonetheless, we are mindful that the current version will always be a work in progress, and both of us welcome your comments on any aspect of it. You can reach us by email at patt@ece.utexas.edu and sjp@crhc.uiuc.edu. We hope you will.

Yale N. Patt Sanjay J. Patel

## **CONTENTS**

Cha	pter 1	2.6	Operations on Bits-Part II: Logical
We	lcome Aboard 1		Operations 29
***	reome risourd 1		2.6.1 The AND Function 29
1.1	What We Will Try To Do 1		2.6.2 The OR Function 30
1.2	How We Will Get There 2		2.6.3 The NOT Function 30
1.3	A Computer System 3		2.6.4 The Exclusive-OR Function 31
1.4	Two Very Important Ideas 5		2.6.5 Examples 31
1.5	Computers as Universal Computational	2.7	Other Representations 32
1.5	Devices 6		2.7.1 Floating Point Data Type 32
1 6	How Do We Get the Electrons to Do the		2.7.2 ASCII Codes 36
1.6	Work? 8		2.7.3 Hexadecimal Notation 36
			Problems 38
	1.6.2 The Algorithm 10	Cha	oter 3
	1.6.3 The Program 10 1.6.4 The ISA 11	Die	ital Logic Structures 47
	1.6.5 The Micoarchitecture 12	פוט	ital Logic Structures 47
	1.6.6 The Logic Circuit 12	3.1	The Transistor 47
	1.6.7 The Devices 13	3.2	Logic Gates 49
	1.6.8 Putting It Together 13	3.2	3.2.1 The NOT Gate (or, Inverter) 49
	Problems 13		3.2.2 OR and NOR Gates 50
	1100101110		3.2.3 AND and NAND Gates 52
Chapter 2			3.2.4 DeMorgan's Law 54
-			3.2.5 Larger Gates 55
Bits	s, Data Types, and Operations 17	3.3	Combinational Logic Structures 55
2.1	Pits and Data Tymes 17	5.5	3.3.1 Decoder 56
2.1	Bits and Data Types 17 2.1.1 The Bit as the Unit of Information 17		3.3.2 Mux 56
			3.3.3 Full Adder Circuit 57
	• 1		3.3.4 Logical Completeness 60
2.2	Integer Data Types 19 2.2.1 Unsigned Integers 19	3.4	
	2.2.1 Onsigned integers 19 2.2.2 Signed Integers 19	3.4	Basic Storage Elements 60 3.4.1 The R-S Latch 61
2.3			3.4.1 The R-S Latch 61 3.4.2 The Gated D Latch 62
	2's Complement Integers 21		
2.4	Binary-Decimal Conversion 23 2.4.1 Binary to Decimal Conversion 23	25	8
		3.5	The Concept of Memory 64
	The state of the s		3.5.1 Address Space 64
2.5	Operations on Bits—Part I: Arithmetic 25		3.5.2 Addressability 64
	2.5.1 Addition and Subtraction 25	2.6	3.5.3 A $2^2$ -by-3-bit Memory 65
	2.5.2 Sign-extension 26	3.6	The Data Path of the LC-2 67
	2.5.3 Overflow 27		Problems 68

#### Contents

Cha	pter 4		5.6	The Data Path Revisited 110
Th	e Von N	Neumann Model 75		5.6.1 Basic Components of the Data Path 110
	0 1011 1	100111111111111111111111111111111111111		5.6.2 The Instruction Cycle 113
4.1	Basic (	Components 75		Problems 114
	4.1.1	Memory 75		
	4.1.2	Processing Unit 77		
	4.1.3	Input and Output 78		
	4.1.4	Control Unit 78		
	4.1.5	Summary: The LC-2 as an Example of		
		the Von Neumann Model 79		
4.2		tion Processing 79		
	4.2.1	The Instruction 80		
	4.2.2	The Instruction Cycle 82		
	4.2.3	Examples 84		
4.3		ng the Sequence of Execution 84		
4.4	10.00	ng the Computer 85		
	Problem	is 87		
Cha	pter 5			
		0.1		
Ine	LC-2	91		
5.1	The ISA	A: Overview 91		
	5.1.1	Memory Organization 92		
	5.1.2	Registers 92		
	5.1.3	The Instruction Set 93		
	5.1.4	Opcodes 93		
	5.1.5	Data Types 93		
	5.1.6	Addressing Modes 95		
	5.1.7	Condition Codes 95		
5.2		Instructions 95		
5.3		ovement Instructions 97		
	5.3.1	Immediate Mode 98		
	5.3.2	Direct Mode 98		
	5.3.3	Indirect Mode 99		
	5.3.4	Base+Offset Mode 99		
	5.3.5	An Example 99		
5.4		Instructions 101		
	5.4.1	Conditional Branches 101		
	5.4.2	An Example 102		
	5.4.3	Two Methods for Loop Control 104		
	5.4.4	Example: Adding a Column of		
	5.4.5	Numbers Using a Sentinel 105 The TRAP Instruction 106		
5.5		Example: Counting Occurrences of		
		eter 107		
	- Ciiui ac	101		

chapter

1

## Welcome Aboard

#### 1.1 WHAT WE WILL TRY TO DO

Welcome to "From Bits and Gates to C and Beyond." Our intent is to introduce you over the next 515 pages to the world of computing. As we do so, we have one objective above all others: to show you very clearly that there is no magic to computing. The computer is a deterministic system—every time we hit it over the head in the same way and in the same place (provided, of course, it was in the same starting condition), we get the same response. The computer is not an electronic genius; on the contrary, if anything, it is an electronic idiot, doing exactly what we tell it to do. It has no mind of its own.

What appears to be a very complex organism is really just a huge, systematically interconnected collection of very simple parts. Our job throughout this book is to introduce you to those very simple parts, and, step-by-step, build the interconnected structure that you know by the name *computer*. Like a house, we will start at the bottom, construct the foundation first, and then go on to add layers and layers, as we get closer and closer to what most people know as a full-blown computer. Each time we add a layer, we will explain what we are doing, tying the new ideas to the underlying fabric. Our goal is that when we are done, you will be able to write programs in a computer language such as C, using the sophisticated features of that language, and understand what is going on underneath, inside the computer.

#### 1.2 HOW WE WILL GET THERE

We will start (in Chapter 2) by noting that the computer is a piece of electronic equipment and, as such, consists of electronic parts interconnected by wires. Every wire in the computer, at every moment in time, is either at a high voltage or a low voltage. We do not differentiate exactly how high. For example, we do not distinguish voltages of 115 volts from voltages of 118 volts. We only care whether there is or whether there is not a large voltage relative to 0 volts. That absence or presence of a large voltage relative to 0 volts is represented as 0 or 1.

We will encode all information as sequences of 0s and 1s. For example, one encoding of the letter a that is commonly used is the sequence 01100001. One encoding of the decimal number 35 is the sequence 00100011. We will see how to perform operations on such encoded information.

Once we are comfortable with information represented as codes made up of 0s and 1s and operations (addition, for example) being performed on these representations, we will begin the process of showing how a computer works. In Chapter 3, we will see how the transistors that make up today's microprocessors work. We will further see how those transistors are combined into larger structures that perform operations, such as addition, and into structures that allow us to save information for later use. In Chapter 4, we will combine these larger structures into the Von Neumann machine, a basic model that describes how a computer works. In Chapter 5, we will begin to study a simple computer, the LC-2. LC-2 stands for Little Computer 2; we started with LC-1 but needed a second chance before we got it right! The LC-2 has all the important characteristics of the microprocessors that you may have already heard of, for example, the Intel 8088, which was used in the first IBM PCs back in 1981. Or the Motorola 68000, which was used in the Macintosh, vintage 1984. Or the Pentium III, one of the high-performance microprocessors of choice in the PC of the year 2000. That is, the LC-2 has all the important characteristics of these "real" microprocessors, without being so complicated that it gets in the way of your understanding.

Once we understand how the LC-2 works, the next step is to program it, first in its own language (Chapter 6), then in a language called *assembly language* that is a little bit easier for humans to work with (Chapter 7). Chapter 8 deals with the problem of getting information into (input) and out of (output) the LC-2. Chapter 9 covers two sophisticated LC-2 mechanisms, TRAPs and subroutines.

We conclude our introduction to programming the LC-2 in Chapter 10 by first introducing two important concepts (stacks and data conversion), and then by showing a sophisticated example: an LC-2 program that carries out the work of a handheld calculator.

In the second half of the book (Chapters 11–19), we turn our attention to a high-level programming language, C. We include many aspects of C that are usually not dealt with in an introductory textbook. In almost all cases, we try to tie high-level C constructs to the underlying LC-2, so that you will understand what you demand of the computer when you use a particular construct in a C program.

Our treatment of C starts with basic topics such as variables and operators (Chapter 12), control structures (Chapter 13), and functions (Chapter 14). We then move on to the more advanced topics of debugging C programs (Chapter 15), recursion (Chapter 16), and pointers and arrays (Chapter 17).

We conclude our introduction to C by examining two very common high-level constructs, input/output in C (Chapter 18) and the linked list (Chapter 19).

#### 1.3 A COMPUTER SYSTEM

We have used the word *computer* many times in the preceding paragraphs, and although we did not say so explicitly, we used it to mean a mechanism that does two things: It directs the processing of information and it performs the actual processing of information. It does both these things in response to a computer program. When we say "directing the processing of information," we mean figuring out which task should get carried out next. When we say "performing the actual processing," we mean doing the actual additions, multiplications, and so forth that are necessary to get the job done. A more precise term for this mechanism is a central processing unit (cpu), or simply a processor. This textbook is primarily about the processor and the programs that are executed by the processor.

Twenty years ago, the processor was constructed out of ten or more 18-inch electronic boards, each containing 50 or more electronic parts known as integrated circuit packages (see Figure 1.1). Today, a processor usually consists of a single microprocessor chip, built on a piece of silicon material, measuring less than an inch square, and containing many millions of transistors (see Figure 1.2).

However, when most people use the word *computer*, they usually mean more than the processor. They usually mean the collection of parts that in combination form their "computer system" (see Figure 1.3). A computer system usually includes, in addition to the processor, a keyboard for typing commands, a mouse for clicking on menu entries, a monitor for displaying information that the computer system has produced, a printer for obtaining paper copies of that information, memory for temporarily storing information, disks and CD-ROMs of one sort or another for storing information for a very long time, even after the computer has been turned off, and the collection of programs (the software) that the user wishes to execute.

These additional items are useful to help the computer user do his/her job. Without a printer, for example, the user would have to copy by hand what is displayed on the monitor. Without a mouse, the user would have to type each command, rather than simply click on the mouse button.

So, as we begin our journey which focuses on how we get less than 1 square inch of silicon to do our bidding, we note that the computer systems we use contain a lot of other components to make our life more comfortable.

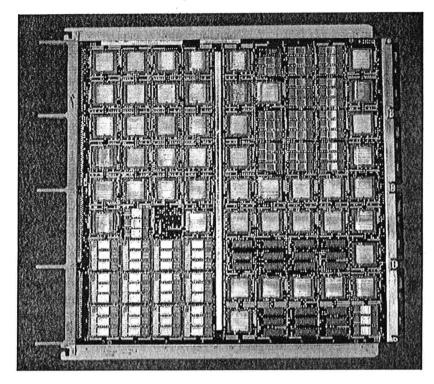


Figure 1.1 A processor board, vintage 1980s. (Courtesy of Emilio Salgueiro, Unisys Corporation.)

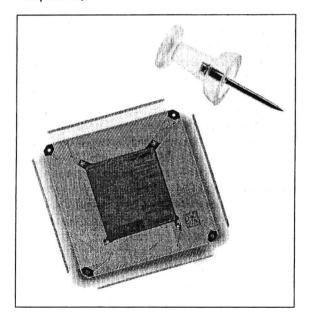


Figure 1.2 A microprocessor, vintage 1998. (Courtesy of Intel Corporation.)

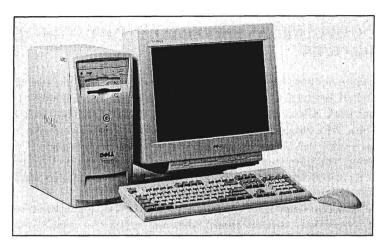


Figure 1.3 A personal computer. (Courtesy of Dell Computer.)

#### 1.4 TWO VERY IMPORTANT IDEAS

Before we leave this first chapter, there are two very important ideas that we would like you to understand, ideas that are at the core of what computing is all about.

IDEA 1: All computers (the biggest and the smallest, the fastest and the slowest, the most expensive and the cheapest) are capable of computing exactly the same things if they are given enough time and enough memory. That is, anything a fast computer can do, a slow computer can do also. The slow computer just does it more slowly. A more expensive computer cannot figure out something that a cheaper computer is unable to figure out as long as the cheap computer can access enough memory. (You may have to go to the store to buy disks whenever it runs out of memory in order to keep increasing memory.) All computers are able to do exactly the same things. Some computers can do things faster, but none can do more than any other.

IDEA 2: We describe our problems in English or some other language spoken by people. Yet the problems are solved by electrons running around inside the computer. It is necessary to transform our problem from the language of humans to the voltages that influence the flow of electrons. This transformation is really a sequence of systematic transformations, developed and improved over the last 50 years, which combine to give the computer the ability to carry out what appears to be some very complicated tasks. In reality, these tasks are simple and straightforward.

The rest of this chapter is devoted to discussing these two ideas.



