

# Discrete Mathematical Structures

for Computer Scientists and Engineers

M.K. Das

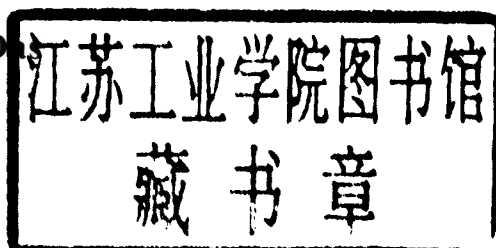


Alpha  
Science

# Discrete Mathematical Structures

for Computer Scientists and Engineers

M.K. D



Alpha Science International Ltd.  
Oxford, U.K.

**M.K. Das**

Institute of Information and Communication  
University of Delhi South Campus  
New Delhi, India

Copyright © 2007

Reprint 2008

Alpha Science International Ltd.  
7200 The Quorum, Oxford Business Park North  
Garsington Road, Oxford OX4 2JZ, U.K.

**[www.alphasci.com](http://www.alphasci.com)**

All rights reserved. No part of this publication may be reproduced, stored in a retrieval system, or transmitted in any form or by any means, electronic, mechanical, photocopying, recording or otherwise, without prior written permission of the publisher.

ISBN 978-1-84265-298-5

Printed in India

# Discrete Mathematical Structures

for Computer Scientists and Engineers

**In Memory of  
Sri Surendra Chandra Das and Smt. Renu Bala Das  
my parents**

# Preface

The present book is the result of teaching the course *IT-14, Mathematical Foundation for Computer Science* at Institute of Informatics and Communication (IIC), University of Delhi South Campus. The present text is suitable for a one- or two-semester course and covers topics incorporated in the course work for students of Computer Science, Computer Application, Electronics & Communication, Telecommunication Engineering and Informatics and Information Technology at Indian and other Universities. I was guided to write this book because of my interest in the subject of *Discrete Mathematics* and *Theoretical Computer Science*. Keeping this in view, the primary emphasis in the book is to discuss basic issues in *Discrete Mathematical Structures* and also in the area of *Mathematical Theory of Computation*. At IIC, we have students coming from different streams such as Mathematical Sciences, Physical Sciences, Electronics and Engineering. Even though they are quite familiar and conversant with basic algebra, calculus, differential equations, integral transforms etc., it is observed that they find the subject of discrete mathematics as somewhat difficult. Further discussion with many of them, during last several years, suggest that the students in general want to focus on the relevance and practicality of various ideas involved in the area of Discrete Mathematics. This aspect has been taken into account during various stages of this project. It is already known that the discrete mathematics is the mathematics of discrete objects and their binding relationships. Developments in digital technology during the second half of the last century coincided with interest in the subject of discrete mathematics. It is now established that the various practical issues involved in the theory of computation involves basic understanding of some abstract structures and their formulations as discussed in the text. The understanding of such formulations further helps in developing ideas involved in data structures and algorithms and their performance. These ideas along with the concept of finite state machine and automata enables one to visualize and formulate the basic ideas of the theory of computation. In the present text I have tried to concisely present various conceptual ideas involved at various level with examples of their applications. The emphasis is made to show the relevance and practical application of the ideas involved in discrete mathematics. I feel that the present text is self contained and user friendly and hopefully the language is simple for students and amateurs to understand and appreciate the basic ideas of mathematics as used in the theory of computation. Suggestions for improvement are welcome and may be sent to [mrinal@iic.ac.in](mailto:mrinal@iic.ac.in)

**M.K. Das**  
IIC, UDSC



# Acknowledgements

I would like to thank Prof. Mahender Singh (retired) of Indian Institute of Technology, Kharagpur for going through the initial draft of various chapters in the book and providing me very useful suggestions. I would like to thank Prof. Dinesh Singh, Director, UDSC for providing necessary support, infrastructure and encouragement at IIC. Dr. Sanjeev Singh my colleague at IIC, UDSC helped and provided very useful inputs at various stages of preparation of the manuscript and I am thankful to him. Many students, notably Mr. Rajat Budhiraja, Ms. Purnima Jain, Mr. Amritabha Mukherjee, Mr. Ashwini Dogra, Ms. Anindita Verma, Mr. Karthik Kamath, Ms. Juhi Garg, Mr. Mohit Dixit, Ms. Priyanka Rana, Mr. Shivam Gupta, Ms. Sonia Bhattacharya, Ms. Geetanjali, Mr. Rashim Pathak, Mr. Parmjeet Singh, Mr. Nikhil Adhikary, Mr. Praveen, Mr. Vipin Mishra, Mr. Pushkar Raj, Ms. Prachi Garg, Madhur Sethi of IIC helped in reading the draft of various chapters and pointing out corrections needed. I owe my thanks to all of them. I thank Prof. R.S. Gupta (OSD, IIC), Dr. Mridula Gupta of Department of Electronic Science, UDSC for providing encouragement at various stages. I am thankful to Dr. A. Sankara Reddy, Principal, Sri Venkateswara College (SVC), University of Delhi for provided me support, encouragement at the initial stages of this work. I thank my colleagues and friends Dr. L.M. Saha, Department of Mathematics, Zakir husain College, New Delhi; Dr. V. Panda, Dr. S.K. Khurana, Dr. V. Chandrasekhar Rao of Department of Chemistry, SVC, Dr. R.K. Pal, Dr. (Mrs.) Renu Jain, Dr. B.V.G. Rao, Dr. Pratima Sharma, Dr. C. Kaur, Dr. B. Biswal, Dr. C. Singh and Dr. A.K. Chaudhary of the Department of Physics, SVC for encouragement, interest and support for the work during my earlier tenure at SVC. Finally I would like to thank authors whose write up, research articles and texts on the subject inspired and fascinated me to the subject of discrete mathematics.

My sincere thanks goes to Dr. Sandeep Kumar, Sponsoring Editor and Mr. N.K. Mehra, Publisher and Managing Director, Narosa Publishing House, New Delhi for their unwavering support and encouragement for this project.

Finally, I am grateful to Soumya Das, my wife, and Sauresh Das, my son for their patience, encouragement during many long days, nights and weekends that were used to complete this project.

I take the responsibility for any error(s) that still remain in this book. I would certainly appreciate receiving any comments and suggestion for further improvements and presentation from readers.

M.K. Das

# Contents

*Preface*

vii

*Acknowledgements*

ix

## **1. Elements of Set Theory**

1

- 1.1 Basic Notation 1
  - 1.1.1 Subsets 2
  - 1.1.2 Equality of Sets 3
- 1.2 Power Set 3
- 1.3 Complement of a Set 3
- 1.4 Universal Set 3
- 1.5 Null Set 3
- 1.6 Basic Set Operations 4
- 1.7 Finite Set 7
- 1.8 Algebraic Properties of Set Operations 7
  - 1.8.1 Principle of Inclusion and Exclusion 8
- 1.9 Characteristic Function 9
  - 1.9.1 Hamming Distance 10
- 1.10 Computer Representation of Sets and Subsets 10
- 1.11 Concept of a Fuzzy Set 11
- 1.12 Definition of a Fuzzy Set 11
- 1.13 Membership Function and Characteristics of Fuzzy Sets 13
- 1.14 Fuzzy Sets and Sets 17
- 1.15 Fuzzy Set Operations 17

## **2. Mathematical Logic**

27

- 2.1 Propositions and Logical Connectives 28
  - 2.1.1 Negation 28
  - 2.1.2 Conjunction 28
  - 2.1.3 Disjunction 29
  - 2.1.4 Conditional Connectives 31
  - 2.1.5 Formulas and Rule of Precedence 33
  - 2.1.6 Tautologies 33
  - 2.1.7 Logical Equivalence 35
  - 2.1.8 Logical Quantifiers 36
- 2.2 Method of Proof 37
- 2.3 Mathematical Induction 40



<b>3. Relations</b>	<b>50</b>
3.1 Cross-Product or Cartesian Product of Two Sets	50
3.2 Binary Relation	51
3.3 N-ary Relations	52
3.4 Matrix of a Relation	53
3.5 Digraphs	54
3.6 Sets Related to R	55
3.7 Paths in Relations	56
3.8 Properties of Relations in a Set	59
3.8.1 Digraph of Reflexive, Symmetric, Asymmetric, Antisymmetric, and Transitive Relations	60
3.8.2 Matrix of Reflexive, Symmetric, Asymmetric, Antisymmetric and Transitive Relations	62
3.9 Equivalence Relation	63
3.10 Partition of a Set	65
3.11 Interpretation of Equivalence Relation using Digraph	65
3.12 Equivalence Relations and Partitions	66
3.13 Operations on Relation	69
3.13.1 Composition of Relations	71
3.13.2 Closures of Relations	73
3.13.3 Warshall Algorithm	77
3.14 Linked List Representation of Data	85
3.15 Computer Representation of Relations and Digraph	88
<b>4. Functions</b>	<b>97</b>
4.1 One-to-one, onto and Inverse Functions	101
4.2 Composition of Functions	107
4.3 The Graphs of Functions	109
4.4 Some Functions in Computer Science	110
4.5 Growth of Functions	120
4.6 Recursive Functions	124
<b>5. Partial Order and Structures</b>	<b>136</b>
5.1 Partially Ordered Sets	136
5.2 Lexicographic Order	139
5.3 Hasse Diagram	141
5.4 Maximal and Minimal Elements of a Poset	145
5.5 Topological Sorting	151
5.6 Lattice	154
5.6.1 Properties of Lattices	157
5.6.2 Lattice as Algebraic System	159
5.6.3 Direct Product of Lattices	162
5.6.4 Sublattice	163
5.6.5 Lattice Homomorphism	165
5.7 Special Lattices	166

5.8	Finite Boolean Algebra	168	
5.8.1	Boolean Functions	172	
5.8.2	Representation of Basic Boolean Polynomials: Logic Gates	175	
5.8.3	Minimization of Combinational Circuits	177	
<b>6.</b>	<b>Combinatorics and Algebraic Systems</b>		<b>198</b>
6.1	Basic Counting Methods	198	
6.2	Permutation and Combination	200	
6.3	Elements of Discrete Probability Theory	203	
6.3.1	Probability Axioms	203	
6.3.2	Conditional Probability	204	
6.3.3	Independence of Events	205	
6.3.4	Baye's Rule	206	
6.3.5	Bernoulli Trials	208	
6.3.6	Random Variables	210	
6.3.7	Distribution of a Random Variable	210	
6.4	Recurrence Relations	212	
6.4.1	Solution of Recurrence Relations	214	
6.4.2	Generating Function and its Application to find Solution of Recurrence Relation	219	
6.4.3	Divide-and-Conquer Relations	222	
6.5	Algebraic Systems	224	
6.5.1	Binary Operation	224	
6.5.2	Semigroup	225	
6.5.3	Monoid	226	
6.5.4	Subsemigroup and Submonoid	227	
6.5.5	Semigroup Isomorphism and Homomorphism	227	
6.5.6	Products and Quotients (or Factor) Semigroups	228	
6.5.7	Groups	231	
6.5.8	Finite Groups and Group Tables	233	
6.5.9	Elementary Properties of Groups	234	
6.6	Permutation Group	236	
6.6.1	Cyclic Groups	239	
6.6.2	Subgroup	241	
6.6.3	Group Homomorphism	241	
6.6.4	Cosets and Lagrange Theorem	242	
6.6.5	Normal Subgroup	244	
6.6.6	Quotient Group and Product of Groups	245	
6.6.7	Algebraic Coding Theory	245	
<b>7.</b>	<b>Elements of Graph Theory</b>		<b>265</b>
7.1	Graphs	265	
7.2	Some Graph Models	269	
7.3	Basic Graph Terminology	272	
7.4	Special Simple Graphs	277	
7.4.1	Application of Simple Special Graphs	279	

7.5	Operations on Graphs	281	
7.6	Graph Representation	287	
7.6.1	Adjacency List	287	
7.6.2	Matrix Representation of a Graph	288	
7.7	Isomorphism of Graphs	291	
7.8	Counting Paths of Length $k$	296	
7.9	Graph Traversal: Eulerian and Hamiltonian Paths	299	
7.9.1	Euler Graph	299	
7.9.2	Fleury's Algorithm	302	
7.9.3	Hamilton Graph	305	
7.9.4	Hamilton Circuit in a $n$ -cube	309	
7.9.5	Conversion of Analog Information to Digital Form	310	
7.10	Shortest Path in a Graph	311	
7.10.1	A Shortest Path Algorithm	311	
7.10.2	The Travelling Salesman Problem	321	
7.11	Planar Graphs	323	
7.12	Detection of Planarity and Kuratowski's Theorem	327	
<b>8.</b>	<b>Trees</b>		<b>342</b>
8.1	Trees	342	
8.2	Tree Terminology	345	
8.3	Rooted Labeled Trees	352	
8.4	Prefix Code	357	
8.5	Binary Search Tree	366	
8.6	Tree Traversal	370	
8.6.1	Prefix, Infix and Postfix Notations	375	
8.7	Spanning Tree	381	
8.7.1	Breadth—First Search Method (BFS)	384	
8.7.2	Depth—First Search Method (DFS)	386	
8.8	Minimum Spanning Trees (MST)	390	
8.9	Decision Trees	397	
8.10	Sorting Methods	399	
8.10.1	The Bubble Sort	400	
8.10.2	The Merge Sort	402	
<b>9.</b>	<b>Finite State Machine and Automata</b>		<b>420</b>
9.1	Finite State Model (FSM)	421	
9.2	Input and Output String in FSM	425	
9.3	Input–Output Transformation of an FSM	429	
9.4	Equivalent Machines	430	
9.5	Finite State Machine and Pattern Recognition (Finite State Automata)	434	
9.5.1	Deterministic Finite Automaton (DFA)	434	
9.5.2	Nondeterministic Finite Automata	439	
9.6	Conversion of <del>NDFA</del> into DFA	441	
9.7	The Equivalence of DFA and NDFA	445	

9.8	Finite Automata with $\Lambda$ -moves	447
9.9	Regular Sets and Regular Expressions	450
9.9.1	Identities for Regular Expressions	452
9.9.2	Regular Expressions (Regular Sets) and Finite Automata	454
9.9.3	Regular Expression as Recognized by a Finite Automata	459
9.10	Regular Language and Regular Expression	462
<b>10.</b>	<b>Languages, Grammar, Push Down Automata &amp; Turing Machine</b>	<b>473</b>
10.1	Languages and Grammar	473
10.2	Types of Grammar (Chomsky's Classification)	479
10.3	Regular Grammar	481
10.3.1	Regular Grammar and Regular Language	481
10.3.2	Closure Properties of Regular Languages	484
10.3.3	Identification of Regular Languages	487
10.4	Context-Free Grammar & Language (or Language with Type-2 Grammar)	490
10.4.1	Ambiguity in Context-free Grammar	492
10.4.2	Removing Ambiguity from CFG	494
10.5	Context-free Grammar and Normal Form	494
10.5.1	Elimination of Useless Symbols	495
10.5.2	Elimination of $\Lambda$ -productions	498
10.5.3	Elimination of Unit Productions	500
10.5.4	Chomsky Normal Form	501
10.5.5	Greibach Normal Form	503
10.6	Pumping Lemma for CFG	505
10.7	Closure Properties of CFL	507
10.8	Push Down Automata	507
10.8.1	Moves in PDA	509
10.8.2	Instantaneous Descriptions	510
10.8.3	Language accepted by PDA's	511
10.9	PDA and Context-free Grammar	514
10.10	Turing Machine	516
10.10.1	Instantaneous Description (ID)	518
10.10.2	Moves in TM	518
10.10.3	Representation of Transition	518
10.10.4	TM Computable Languages	520
10.10.5	TM Computable Functions	520
	<i>References</i>	529
	<i>Index</i>	532

# 1

---

## Elements of Set Theory

A subject called set theory has been identified as even more basic than arithmetic and geometry by mathematicians during the last century. It provides important basic concepts of contemporary mathematics. The ideas of set theory find application in all branches of mathematics as they form a powerful language for reasoning about various mathematical objects. Intuitively we define a *set as collection of objects*. The objects in such a collection are referred to as *set elements*. Even though such a definition lacks rigors, we use it in much the same way as a point and a line are left undefined in geometry. In this chapter, we assume that a set is determined when a plurality (consisting of a number of things) is bunched together into a single entity. For instance a set called Hindi alphabet comprises 49 things called letters. Similarly the real number system could be considered as a set formed collectively from things each of which is a real number.

### 1.1 Basic Notation

In order to make the concept of a set more precise, we distinguish things which are collected together in a set from those which are not. Because of this, the set notation and theory must be precise and unambiguous. It is customary to designate a set by capital letters A, B, P, Q, X, Y etc. When the number of things or objects comprising a set is small, the set is denoted by listing the elements within the braces {}, thus establishing the singular nature of the set. For instance a set of digits in the binary system of enumeration is denoted by  $X = \{0, 1\}$ . Similarly  $A = \{a, e, i, o, u\}$  denotes the set of vowels in English language.

We represent the relationship between a set and its elements by the symbol  $\in$  (belongs to) in the following way:

$$i \in V$$

means that the element  $i$  is a member of the set named  $V$ . Alternatively such a representation also means ' $i$  belongs to  $V$ ' or ' $i$  is in  $V$ ' or ' $i$  is an element of  $V$ '. If however, an element say  $d$  does not belong to a set  $V$  above, we write

$$d \notin V.$$

For a large set, the enumeration could be done without actually listing all the elements. For instance, 26 letters of the set English alphabet could be described as:

$$A = \{a, b, c, \dots, x, y, z\}$$

where three consecutive dots, called ellipsis, mean that the list should be continued in accordance with the pattern indicated. For instance the set of natural numbers could be represented as:

$$N = \{1, 2, 3, \dots\}.$$

Similarly the set of integers is represented as:

$$Z = \{\dots, -2, -1, 0, 1, 2, \dots\}$$

Another but convenient notation of writing a set is as:

$$B = \{x: P(x)\},$$

where  $x$  is a variable and  $P(x)$  is a descriptive phrase called *predicate*. The variable is mostly indicated by a lower case letter and the predicate is a description of the allowed replacement for that variable. They are usually separated by a colon : or by a vertical bar |. For instance the set  $B = \{x: x \text{ is an integer, } x < 0\}$ , means that  $B$  is a set of  $x$  such that  $x$  is an integer and  $x$  is less than zero. Similarly the set  $E = \{x: x^2 - 4 = 0\}$  means that the set  $E$  has elements  $x$  such that  $x^2 - 4 = 0$ . Since  $x = \pm 2$  are the solutions of the equation  $x^2 - 4 = 0$ , we have  $E = \{2, -2\}$ . The two forms are equivalent.

Further we would like to mention here that the order in which the elements of a set are listed is not important. Thus the representation of the set  $\{1, 5, 2\}$  by  $\{1, 2, 5\}$ ,  $(2, 5, 1)$ ,  $\{5, 1, 2\}$  etc., are all equivalent. Moreover the listing in a set might have repeated element(s) and this may be ignored. Thus  $\{1, 5, 2, 5, 1\}$  is another representation of  $\{1, 5, 2\}$ .

### 1.1.1 Subsets

We may write the following definition for the subset of a set:

**Definition:** Let  $A$  and  $B$  be two sets. If every element of  $A$  is also an element of  $B$ , then  $A$  is a subset of  $B$  and is written as  $A \subseteq B$ .

Here the symbol  $\subseteq$  is called the *inclusion* or *subset relation*. Thus  $A \subseteq B$  also means ' $A$  is included in  $B$ ' or ' $A$  is contained in  $B$ '. If  $A$  is not a subset of  $B$ , we write  $A \not\subseteq B$ . For instance if  $A = \{2, 5, 7\}$  and  $B = \{5, 7, 2, 3, 9\}$ , we find  $A \subseteq B$  since for every  $x \in A$ , we have  $x \in B$ . Similarly if  $N, Z, Q$  and  $R$  respectively denote the set of positive integer, set of integers, set of rational numbers and set of real numbers then

$$N \subseteq Z \subseteq Q \subseteq R.$$

In the foregoing definition on subset, observe that the possibility of  $A = B$  is not excluded. Excluding such a possibility reduces the set as proper subset of  $B$  or  $A \subset B$ .

Also if  $A$  be a set and  $B = \{A, \{A\}\}$ . We find  $A \in B$  and  $\{A\} \in B$  and  $\{\{A\}\} \subseteq B$ . However  $A \not\subseteq B$ .



**Example 1** If  $A = \{1, 2, 3\}$  and  $B = \{\{1, 2, 3\}, 1, 2, 3\}$ . What can be said about the set relationship between  $A$  and  $B$ ?

**Solution** It is readily seen that  $A \subseteq B$  and  $A \in B$ .

### 1.1.2 Equality of Sets

We have the following definition for the equality of sets:

**Definition:** Let  $A$  and  $B$  be two sets. Then two sets are equal i.e.,  $A = B$ , if and only if for every  $x \in A$ ,  $x \in B$  and vice-versa. In other words if  $A \subseteq B$  and  $B \subseteq A$  then  $A = B$ .

**Example 2** Consider two sets  $A = \{2, 1, 3, 6\}$  and  $B = \{6, 1, 2, 3\}$ . Since every element of  $A$  is also an element of  $B$  and vice-versa, we have  $A = B$ . Note that the ordering of elements of  $A$  or  $B$  is unimportant.

## 1.2 Power Set

We define the power set as follows:

**Definition:** Let  $A$  be a set consisting of  $n$  elements. The set of all subsets of  $A$  is called the power set of  $A$ . It consists of  $2^n$  elements.

**Example 3** For the set  $A = \{a, b, c, d\}$ , find the power set  $\mathcal{P}(A)$  of  $A$ ?

**Solution** Here we find four element subset as  $A = \{a, b, c, d\}$ ; three element subsets as  $\{a, b, c\}, \{b, c, d\}, \{a, c, d\}, \{a, b, d\}$ ; two element subsets as  $\{a, b\}, \{a, c\}, \{a, d\}, \{b, c\}, \{b, d\}, \{c, d\}$ ; and one element subsets as  $\{a\}, \{b\}, \{c\}, \{d\}$ . Further zero element subset is  $\phi$ . Thus in all there are 16 subsets or  $2^4$  subsets of  $A$ . The power set  $\mathcal{P}(A)$  is also denoted as  $2^A$ .

## 1.3 Complement of a Set

**Definition:** Let  $A$  and  $B$  be two sets. The set consisting of all elements of  $B$  which are not in  $A$  is called the complement of  $A$  in  $B$  and is denoted by  $B - A$ .

In the set builder notation, we may write  $B - A = \{x: x \in B, x \notin A\}$ . To find the complement of  $A$  in  $B$ , we eliminate in  $B$  all elements which belong to  $A$ .

For instance  $\mathbb{Z} - \mathbb{N}$  corresponds to a set containing the element zero and all negative integers.

## 1.4 Universal Set

In a set the elements usually belong to some large set called the universal set  $U$ . A universal set, for instance, in plane geometry consists of all points in the plane; the elements of the set  $A = \{1, 2, 3\}$  may be treated as belonging to a universal set  $U = \mathbb{N}$ .

## 1.5 Null Set

We define a null set or an empty set  $\phi$  as a set that has no element in it. It is also represented as  $\{\}$ . Thus for any set  $A$ , we have  $\phi \subseteq A$  as there are no elements of  $\phi$  that are not in  $A$ . For

instance the square of a real number is always positive. Hence  $\{x: x \text{ is a real number and } x^2 = -3\}$  is a null set  $\phi$ .

## 1.6 Basic Set Operations

Addition and multiplications are examples of binary operations on numbers. These operations could be defined in many other contexts and are not restricted to numbers only. The union, intersection, and complement are some basic operations on sets. Based on these operations, set theory offers a well defined mathematical structure and becomes an indispensable tool in solving quite complicated combinatorial problems.

The basic idea, of a Venn diagram, could be used to understand and visualize the basic set operations and set relationships. In Venn diagram, the largest set under consideration,  $U$ —the universal set is portrayed by a rectangle while its various subsets are represented by circular regions inside it. The Venn diagram (Fig. 1.1) shows the set  $A$  separating  $U$  into two disjoint pieces  $A$  and  $\bar{A}$  (or  $A^c$ )—the complement of  $A$  (shaded region). In the following section, we will define the basic operation on sets and their representation in Venn diagram.

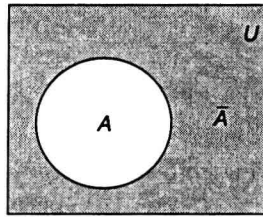


Fig. 1.1 Venn diagram

*Definition:* Let  $A$  and  $B$  be two sets. The set of all elements  $\in A$  or  $\in B$  (or in both  $A$  and  $B$ ) is called the union of  $A$  and  $B$  and is abbreviated as  $A \cup B$ .

Alternatively  $A \cup B = \{x: x \in A \text{ or } x \in B\}$ . The Venn diagram of union of  $A$  and  $B$  is shown below (Fig. 1.2).

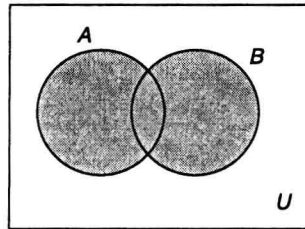


Fig. 1.2 The shaded region showing  $A \cup B$

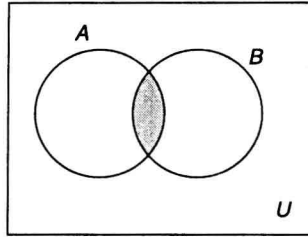
The union of  $n$  sets  $A_i$ 's,  $i = 1, n$  written as:

$$A_1 \cup A_2 \cup A_3 \cup \dots \cup A_n = \bigcup_{i=1}^n A_i$$

represents the set of elements that belongs to one or more of the sets  $A_i$ .

**Definition:** Let  $A$  and  $B$  be two sets. The set of all elements which are both in  $A$  and  $B$  defines the intersection of set  $A$  and  $B$  and is symbolically written as  $A \cap B$ .

Alternatively  $A \cap B = \{x: x \in A \text{ and } x \in B\}$ . The Venn diagram of intersection of two sets  $A$  and  $B$  is shown below (Fig. 1.3).



**Fig. 1.3** The shaded region  $A \cap B$

The intersection of  $n$  sets  $A_i$ 's,  $i = 1, n$  written as:

$$A_1 \cap A_2 \cap A_3 \cap \dots \cap A_n = \bigcap_{i=1}^n A_i$$

represents the set of elements which belongs to all of the sets  $A_i$ .

The following theorem provides the relation between the union and intersection:

**Theorem:** Let  $A$  and  $B$  be two sets. If  $U$  is the universal set then  $\phi \subseteq A \cap B \subseteq A \subseteq A \cup B \subseteq U$ .

**Proof:** The null set is a subset of any set. If  $x \in A \cap B$  then by definition,  $x \in A$  and  $x \in B$ . If  $x \in A$ , then it is in at least one of the sets  $A$  and  $B$ . Hence  $x \in A \cup B$ . Since  $A$  and  $B \in U$ ,  $A \cup B \subseteq U$ .

**Example 5** Let  $A = \{1, 2, 3, 4, 5\}$ ,  $B = \{2, 4, 5, 6, 7\}$  and  $C = \{1, 15, 17\}$ . Find  $A \cup B$ ,  $A \cap B$ ,  $A \cap C$  and  $B \cap C$ ?

**Soution** Using the definition of basic set operations, we find that

$$\begin{aligned} A \cup B &= \{1, 2, 3, 4, 5, 6, 7\} = B \cup A, \\ A \cap B &= \{2, 4, 5\} = B \cap A, \\ A \cap C &= \{1\}, \end{aligned}$$

$B \cap C = \{\} = \phi$  {If two sets are such that no element is common then the two sets are said to be disjoint or nonintersecting}.

**Example 6** If  $A = \{2, 4, 6, 8, 10, \dots\}$ ,  $B = \{1, 3, 5, 7, 9, \dots\}$  then  $A \cup B = N$  and  $A \cap B = \phi$ .

**Example 7** For any set  $A$ ,  $A \cup A = A$  and  $A \cap A = A$ . Therefore, using the foregoing theorem, it is readily seen that  $\phi \in A$ .

**Example 8** The Venn diagram for union and intersection of any three sets could be drawn as (Fig. 1.4).

The complement of set  $A$  with respect to a set  $B$  and vice-versa, defined earlier could be represented by the following Venn diagrams (Fig. 1.5).