

刘彦佩选集

(Selected Publications of Y.P.Liu)

第十四编

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当代科学技术基础理论与前沿问题研究丛书

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Topological Theory on Graphs

图的拓扑理论

刘彦佩 著

中国科学技术大学出版社

· 6371 ·

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内 容 提 要

本书不在于图的拓扑性质本身,而是着意以图为代表的一些组合构形为出发点,揭示与拓扑学中一些典型对象,如多面形、曲面、嵌入、纽结等的联系,特别是显示了定理有效化的途径对于以拓扑学为代表的基础数学的作用。同时,也提出了一些新的曲面模型,为超大规模集成电路的布线尝试构建多方面的理论基础。

本书可作为基础数学、应用数学、系统科学、计算机科学等专业高年级本科生和研究生的补充教材,也可供相关专业的教师和科研工作者参考。

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Preface

The subject of this book reflects new developments mainly by the author himself in company with cooperators most of them his former and present graduate students on the foundation established in Liu, Y.P.[33–34]. The central idea is to extract suitable parts of a topological object such as a graph not necessary to be with symmetry, as linear spaces which are all with symmetry for exploiting global properties in construction of the object. This is a way of combinatorizations and further algebraications of an object via relationship among their subspaces.

Graphs are dealt with three vector spaces over $GF(2)$, the finite field of order 2, generated by 0(dimensional)-cells, 1(dimensional)-cells and 2(dimensional)-cells. The first two spaces were known from, *e.g.*, Lefschetz, S.[2] by taking 0-cells and 1-cells as, respectively, vertices and edges. Of course, a graph is only a 1-complex without 2-cells.

Since the fifties of last century, in Wu, W.J.[1] and Tutte, W.T.[4, 16], the chain groups generated by 0-cells and 1-cells over, respectively, $GF(2)$ and the real field were independently used for describing a graph. And they both then after ten years adopted nonadjacent pair of edges as a 2-cell for which the cohomology on a graph was allowed to be established.

Their results especially in Wu, W.J.[1–6] enabled the present author to create a number of types of planarity auxiliary graphs induced from the graph considered for the study of the efficiency of theorems in Liu, Y.P.[1,2,19,22,42] as one approach. Another approach can be seen in Liu, Y.P.[23–25,43].

More interestingly, two decades later than Liu, Y.P.[1], in Archdeacon, D. and J. Siran[1] a theta graph(network) was used for charac-

terizing the planarity of a given graph. The theta graph can be seen to be a type of planarity auxiliary graph(network) because our planarity auxiliary graphs are subgraphs of the theta graph. However, in virtue of the order of theta network upper bounded by an exponential function of the size of given graph and that of planarity auxiliary network by a quadratic polynomial of the size of given graph, theorems deduced from a theta network are all without efficiency while those from a planarity auxiliary network are all with efficiency.

The effects of planarity auxiliary graphs are reflected in Chapters 8, 10, 11, 12 and 13 with a number of extensions.

On the other hand, in Liu, Y.P.[31] a graph was dealt with a set of polyhedra via double covering the edge set by travels under certain condition so that travels were treated as 2-cells. These enable us to discover homology and another type of cohomology for showing the sufficiency of Eulerian necessary condition in this circumstance. Further, all the results for the planarity of a graph in Whitney, H.[7] on the duality, MacLane, S.[1-2] on a circuit basis and Lefschetz, S.[1] on a circuit double covering have a universal view in this way. In fact, our polyhedra are all on such surfaces, *i.e.*, 2-dimensional compact manifolds without boundary. If a boundary is allowed on a surface, the Eulerian necessary condition is not always sufficient in general. Some person used to have missing the boundary condition in Abrams, L. and C.D. Slilaty[1].

The effects of this theoretical thinking are reflected in Chapters 4,5,7 and 14.

Because of the clarification of the joint tree model of a polyhedron in Liu, Y.P.[35-36] by the present author recently on the basis of Liu, Y.P.[8-9], we are allowed to write a chapter for brief description of the theories of surfaces and polyhedra each in Chapters 2 and 3 with related topics in Chapters 6, 9 and 15.

Although quotient embeddings(current graph and its dual, voltage graph) were quite active in constructing an embedding of a graph on a surface with its genus minimum in a period of decades, this book has no space for them. One reason is that some books have mentioned them such as in White, A.T.[1], Ringel, G.[3] and Liu, Y.P.[33-34], *etc.* Another reason is that only graphs with higher symmetry are suitable for quotient embeddings, or for employing the covering space method whence this book is for general graphs without such a limitation of

symmetry.

In spite of refinements and simplifications for known results, this book still contains a number of new results such as in §5.2, the sufficiency in the proof of Theorem 5.2.1, §9.4, §11.3–4, §13.1–2, §13.4–5 *etc.*, only name a few. Researches were partially supported by NNSF in China under Grants No.60373030 and No.10571013.

Y. P. Liu

Beijing

December 2007

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Chapter 1

Preliminaries

Throughout for the sake of brevity, the usual logical conventions are adopted: disjunction, conjunction, negation, implication, equivalence, universal quantification and existential quantification denoted, respectively, by the familiar symbols: \vee , \wedge , \neg , \Rightarrow , \Leftrightarrow , \forall and \exists . And, $\S x.y$ is for the section y in Chapter x .

In the context, $(i.j.k)$ refers to item k of section j in chapter i .

A reference $[k]$ refers to item k of the corresponding author(s) in the bibliography where k is a positive integer to distinguish publications of the same author(s).

1.1 Sets and relations

A *set* is a collection of objects with some common property which might be numbers, points, symbols, letters or whatever even sets except itself to avoid paradoxes. The objects are said to be *elements* of the set. We always denote elements by italic lower letters and sets by capital ones. The statement “ x is (is not) an element of M ” is written as $x \in M$ ($x \notin M$). A set is often characterized by a property. For example

$$M = \{x \mid x \leq 4, \text{ positive integer} \} = \{1, 2, 3, 4\}.$$

The *cardinality* of a set M (or the number of elements of M if finite) is denoted by $|M|$.

Let A, B be two sets. If $(\forall a) (a \in A \Rightarrow a \in B)$, then A is said to be a *subset* of B which is denoted by $A \subseteq B$. Further, we may define the three main operations: union, intersection and subtraction

respectively as $A \cup B = \{x \mid (x \in A) \vee (x \in B)\}$, $A \cap B = \{x \mid (x \in A) \wedge (x \in B)\}$ and $A \setminus B = \{x \mid (x \in A) \wedge (x \notin B)\}$.

If $B \subseteq A$, then $A \setminus B = A - B$ is denoted by $\overline{B}(A)$ which is said to be the *complement* of B in A . If all the sets are considered as subsets of Ω , then the complement of A in Ω is simply denoted by \overline{A} . The *empty* denoted by \emptyset is the set without element. For those operations on subsets of Ω mentioned above, we have the following laws.

Idempotent law $\forall A \subseteq \Omega, A \cap A = A \cup A = A$.

Commutative law $\forall A, B \subseteq \Omega, A \cup B = B \cup A; A \cap B = B \cap A$.

Associative law $\forall A, B, C \subseteq \Omega, A \cup (B \cap C) = (A \cup B) \cap C$;
 $A \cap (B \cup C) = (A \cap B) \cup C$.

Absorption law $\forall A, B \subseteq \Omega, A \cap (A \cup B) = A \cup (A \cap B) = A$.

Distributive law $\forall A, B, C \subseteq \Omega, A \cup (B \cap C) = (A \cup B) \cap (A \cup C)$;
 $A \cap (B \cup C) = (A \cap B) \cup (A \cap C)$.

Universal bound law $\forall A \subseteq \Omega, \emptyset \cap A = \emptyset, \emptyset \cup A = A; \Omega \cap A = A, \Omega \cup A = \Omega$.

Unary complement law $\forall A \subseteq \Omega, A \cap \overline{A} = \emptyset; A \cup \overline{A} = \Omega$.

The unary complement law is also called the *excluded middle law* in logic.

From the laws described above, we may obtain a large number of important results. Here, only a few is listed for the usage in the context.

Theorem 1.1.1 $\forall A \subseteq \Omega,$

$$\begin{cases} (\forall X \subseteq \Omega) ((A \cap X = A) \vee (A \cup X = X)) \\ \quad \Rightarrow A = \emptyset; \\ (\forall X \subseteq \Omega) ((A \cap X = X) \vee (A \cup X = A)) \\ \quad \Rightarrow A = \Omega. \end{cases} \quad (1.1.1)$$

Theorem 1.1.2 $\forall A, B \subseteq \Omega,$

$$A \cap B = A \Leftrightarrow A \cup B = B. \quad (1.1.2)$$

Theorem 1.1.3 $\forall A, B, C \subseteq \Omega,$

$$(A \cap B = A \cap C) \wedge (A \cup B = A \cup C) \Leftrightarrow B = C. \quad (1.1.3)$$

Theorem 1.1.4 $\forall A \subseteq \Omega$,

$$\overline{\overline{A}} = A. \quad (1.1.4.)$$

Theorem 1.1.5 $\forall A, B \subseteq \Omega$,

$$\overline{A \cup B} = \overline{A} \cap \overline{B}; \overline{A \cap B} = \overline{A} \cup \overline{B}. \quad (1.1.5)$$

From those described above, it is seen that $\overline{\emptyset} = \Omega$ and $\overline{\Omega} = \emptyset$. Further, the symmetry (or duality) that any proposition related to $\cup, \cap, \emptyset, \Omega$ can be transformed into another by interchanging \cup and \cap, \emptyset and Ω .

For $A, B \subseteq \Omega$, an *injection* (or 1 - to - 1 correspondence) between A and B is a mapping $\alpha : A \rightarrow B$ such that $\forall a, b \in A, a \neq b \Rightarrow \alpha(a) \neq \alpha(b)$. A *surjection* between A and B is a mapping $\beta : A \rightarrow B$ such that $(\forall b \in B)(\exists a \in A)(\beta(a) = b)$. If a mapping is both an injection and a surjection, then it is called a *bijection*. Two sets are said to be *isomorphic* if there is a bijection between them. Two isomorphic sets A and B , or write $A \sim B$, are always treated as the same. Of course, for finite sets, it is trivial to justify if two sets are isomorphic by the fact: $\forall A, B \subseteq \Omega, A \sim B \Leftrightarrow |A| = |B|$.

For a set M , let $M \times M = \{ \langle x, y \rangle \mid \forall x, y \in M \}$ which is said to be the *Cartesian product* of M . Here, $\langle x, y \rangle \neq \langle y, x \rangle$ in general.

A *binary relation* R on M is a subset of $M \times M$. The adjective “binary” of the relation will often be omitted in the context. If the relation R holds for $x, y \in M$, then we write $\langle x, y \rangle \in R$, or xRy . An *order*, denoted by \preceq , is a relation R which satisfies the following three laws:

Reflective law $\forall x \in M, xRx$.

Antisymmetry law $\forall x, y \in M, xRy \wedge yRx \Rightarrow x = y$.

Transitive law $\forall x, y, z \in M, xRy \wedge yRz \Rightarrow xRz$.

The set M with the order \preceq is said to be a *poset* (or partial order set) denoted by (M, \preceq) .

Theorem 1.1.6 In a poset (M, \preceq) , $\forall x_1, x_2, \dots, x_n \in M$,

$$x_1 \preceq x_2 \preceq \dots \preceq x_n \preceq x_1 \Rightarrow x_1 = x_2 = \dots = x_n. \quad (1.1.6)$$

The theorem is sometimes called the *anti-circularity law*. If a relation only satisfies Reflective law and Transitive law but not Anti-symmetry law, then it is called the *quasi-order* which is denoted by $\bullet \prec$. A set M with $\bullet \prec$ is said to be a *quoset* denoted by $(M, \bullet \prec)$.

Theorem 1.1.7 Any subset S of a quoset $(M, \bullet \prec)$, is itself a quoset with the restriction of the quasi-order to S .

If a quasi-order R on M satisfies the symmetry law described below, then it is called an *equivalent relation*, or simply an *equivalence* denoted by \sim .

Symmetry law $\forall x, y \in M, xRy \Rightarrow yRx$.

For the equivalence \sim on M , we are allowed to define the set $x(M) = \{y \mid \forall y \in M, y \sim x\}$ which is said to be the *equivalent class* for $x \in M$. The set which consists of all the equivalent classes is called the *quotient set* of (M, \sim) denoted by M/\sim . In a quoset $(M, \bullet \prec)$, let $\sim_{\bullet \prec}$ be defined by

$$\forall x, y \in M, x \sim_{\bullet \prec} y \Leftrightarrow (x \bullet \prec y) \wedge (y \bullet \prec x). \quad (1.1.7)$$

Then, it is easily seen that $\sim_{\bullet \prec}$ is an equivalence on M and that $(M/\sim_{\bullet \prec}, \bullet \prec)$ is also a quoset.

Theorem 1.1.8 A quoset $(M, \bullet \prec)$ is a poset if, and only if, $M/\sim_{\bullet \prec} = M$, or say, it satisfies the anti-circularity law.

In a poset (M, \preceq) , we define the *strict inclusion*, denoted by \prec , of the order by the *anti-reflective law*: $\neg x \in M, x \prec x$ and the *transitive law*: $(x \prec y) \wedge (y \prec z) \Rightarrow x \prec z$ while noticing that $x \preceq y \Leftrightarrow (x \prec y) \vee (x = y)$. If an order \preceq on M satisfies the alternative law described below, then it is called a *total order*, or a *linear order*.

Alternative law $\forall x, y \in M, x \not\preceq y \Rightarrow y \preceq x$.

A set with a total order is said to be a *chain*. The *length* of a chain with n elements is defined to be $n - 1$. From Theorem 1.1.7 and the definitions, we may have

Theorem 1.1.9 Any subset of a poset is a poset and any subset