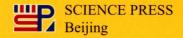
# A Study on Eigenvalues of Higher-Order Tensors and Related Polynomial Optimization Problems

(高阶张量特征值和相关多项式优化问题研究)

Yang Yuning Yang Qingzhi (杨宇宁 杨庆之)



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To the first author's parents:
Yang Cunliang and Huang Linlin
And wife:
Chen Yiying

### **Preface**

In this book, we study the eigenvalues problems of higher order tensors and several specific polynomial optimization problems.

The concepts of eigenvalues and eigenvectors of higher order tensors were presented in 2005 by Qi and Lim independently. Since then, more and more scholars have devoted to this new research field and have made great progress in the past decade. Particularly the various properties of H-eigenvalues of nonnegative tensors and solving methods have been developed. In the first part of this book we systematically introduce related contents. Problems on finding the maximal or minimal H-eigenvalues and Z-eigenvalues are special class of polynomial optimization problems, while polynomial optimization is attracting great attention due to its importance and has been witnessed rapid development in the past years. It is well known that in general a polynomial optimization problem is nonconvex and NP-hard. So people attempt to use suitable relaxed convex optimization problem to approximate a concrete nonconvex optimization problem. Semidefinite programming (SDP) is an important branch of convex optimization and has been studied extensively in theory as well as in algorithms in the past two decades. Besides it is important in itself right, SDP has been found to be very useful as the convex approximation of nonconvex quadratic programming and some discrete optimization problems. In the second part of this book we devote to studying several special polynomial optimization problems by using SDP relaxation strategy, such as maximal real tensor eigenvalue problem and their extended problems.

Most of the contents of this book are composed of PhD dissertation of the first author, which have been published in journals.

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Yang Qingzhi Yang Yuning Nov. 2014

# **Notations**

$x, y, \cdots$	Vectors
$A, B, \cdots$	Matrices
$\mathcal{A},\mathcal{B},\cdots$	Tensors
$\mathbb{R}, \mathbb{C}, \mathbb{Q}$	Real, complex and rational number field
$\mathbb{R}^n, \mathbb{R}^n_+, \mathbb{R}^n_{++}$	n-dimensional real (nonnegative, positive) vector space
$\mathbb{R}^{n \times n}, \mathbb{R}^{n \times n}_+, \mathbb{R}^{n \times n}_{++}$	$n \times n$ dimensional real (nonnegative, positive) matrix space
$S^{n\times n}, S^{n\times n}_+, S^{n\times n}_{++}$	$n \times n$ real symmetric (positive semidefinite, positive definite)
	matrix space
•	$\ell_2$ norm of a vector
$\ \cdot\ _F$	Frobenius norm of a matrix or a tensor
${f T}$	Vector or matrix transpose
$\mathrm{rank}(\cdot)$	Rank of a matrix
$\operatorname{trace}(\cdot)$	Trace of a marix or a tensor
•	Matrix inner product
0	Hadamard operator
$\otimes$	Kronecker product
$\times_d$	Tensor-matrix mode- $d$ product
•	Entry-wise absolute values of a vector, a matrix or a tensor
$\operatorname{diag}(X)$	The diagonal entries of a matrix $X$
Diag(x)	A matrix whose diagonal entries are the vector $x$
$oldsymbol{M}(\cdot)$	Tensor-matrix unfolding operator
$oldsymbol{T}(\cdot)$	Matrix-tensor folding operator
$oldsymbol{V}(\cdot)$	Matrix-vector unfolding operator
$oldsymbol{L}(\cdot)$	Vector-matrix folding operator

# Contents

Chapt	ter 1	Introduction · · · · · · · · · · · · · · · · · · ·	
1.1		nvalues problems of higher order tensors $\cdots 1$	
1.2	Related polynomial optimization problems · · · · · · · · · · · · · · · · · · ·		
1.3	Applications · · · · · · · · · · · · · · · · · · ·		
1.4	Spec	tral properties and algorithms: a literature review $\cdots \cdots 9$	
1.5	The	organization of this book $\cdots 14$	
Chapt	ter 2	Spectral Properties of H-eigenvalue Problems of a	
		Nonnegative Square Tensor · · · · · · 17	
2.1	Intro	oduction · · · · · · 17	
2.2	2.2 From nonnegative matrices to nonnegative tensors · · · · · · · · · · · · · · · · · · ·		
2.3	2.3 Nonnegative irreducible tensors and primitive tensors · · · · · · · · · · · · · · · · · · ·		
2.4	2.4 Perron-Frobenius theorem for nonnegative tensors and related results $\cdots$		
2.5	Geor	metric simplicity · · · · · · · · 29	
2.6	The	Collatz-Wielandt formula · · · · · · · 36	
2.7	Othe	er related results·······41	
2.8	Some properties for nonnegative weakly irreducible tensors		
	2.8.1	Weak irreducibility · · · · · · 46	
	2.8.2	Generalization from nonnegative irreducible tensors to nonnegative	
		weakly irreducible tensors $\cdots 50$	
Chapter 3		Algorithms for Finding the Largest H-eigenvalue of a	
		Nonnegative Square Tensor · · · · · 54	
3.1	Intro	duction · · · · · · · 54	
3.2	-	ly nomial-time approach for computing the spectral radius $\cdots \cdots 55$	
3.3 Two		algorithms and convergence analysis $\cdots \cdots 57$	
	3.3.1	An inexact power-type algorithm $\cdots 58$	
	3.3.2	A one-step inner iteration power-type algorithm $\cdots \cdots 63$	
3.4	Num	erical experiments······65	
	3.4.1	Experiments on the polynomial-time approach $\cdots \cdots 65$	
	3.4.2	Experiments on the inexact algorithms · · · · · · · · · · · · · · · · · · ·	

C	napt	er 4	Spectral Properties and Algorithms of H-singular Value
			Problems of a Nonnegative Rectangular Tensor · · · · · · · 70
	4.1	Intro	oduction · · · · · · · · · · · · · · · · · · ·
	4.2	Preli	minaries 70
4.3 Some			e conclusions concerning the singular value of a nonnegative
recta		recta	angular tensor · · · · · · · · · · · · · · · · · · ·
	4.4 Primitivity and the convergence of the CQZ method for		nitivity and the convergence of the CQZ method for finding the
larg			est singular value of a nonnegative rectangular tensor · · · · · · · · 81
4.5 Algo		Algo	rithms for computing the largest singular value of a nonnegative
		recta	angular tensor·····84
		4.5.1	A polynomial-time algorithm · · · · · · · · · · · · · · · · · · ·
		4.5.2	An inexact algorithm · · · · · · · 85
	4.6	A so	lving method of the largest singular value based on the symmetric
		embe	edding87
		4.6.1	Singular values of a rectangular tensor · · · · · · · 87
		4.6.2	Singular values of a general tensor · · · · · · · 89
C	hapt	er 5	Properties and Algorithms of Z-eigenvalue Problems of
			a Symmetric Tensor·····94
	5.1		oduction·····94
	5.2	Some	e spectral properties·····95
		5.2.1	The Collatz-Wielandt formula · · · · · 95
		5.2.2	Bounds on the Z-spectral radius · · · · · · 99
	5.3	The	reformulation problem and the no duality gap result · · · · · · · · · 100
		5.3.1	The reformulation problem · · · · · · · · · · · · · · · · · · ·
		5.3.2	Dual problem of (RP)
		5.3.3	No duality gap result · · · · · · · · · · · · · · · · · · ·
	5.4	Rela	xations and algorithms · · · · · · · 106
		5.4.1	Nuclear norm regularized convex relaxation of (RP) and the proximal
			augmented Lagrangian method······106
		5.4.2	The truncated nuclear norm regularization and the approximation $\cdots$ 112
		5.4.3	Alternating least eigenvalue method for finding a global minima · · · · · · 114
	5.5	Num	erical results······119
Chapter 6		er 6	Solving Biquadratic Optimization Problems via
			Semidefinite Relaxation · · · · · · 126
	6.1		oduction·····126
6.2 Semio		Semi	definite relaxations and approximate bounds · · · · · · · 127

	6.2.1	The nonnegative case · · · · · · · · · · · · · · · · · · ·	
	6.2.2	The square-free case and the positive semidefinite case $\cdots 130$	
6.3 Appr		oximation algorithms for the biquadratic optimization problems $\cdot \cdot 133$	
	6.3.1	Approximation method for the nonnegative case · · · · · · · · · · · · · · · · · · ·	
	6.3.2	The binary biquadratic optimization problem · · · · · · · · · · · · · · · · · · ·	
	6.3.3	A generalization of the binary biquadratic optimization · · · · · · · · 141	
6.4	Num	erical experiments · · · · · · · · · · · · · · · · · · ·	
Chapter 7		Approximation Algorithms for Trilinear Optimization	
		with Nonconvex Constraints and Extensions · · · · · · · · · 146	
7.1	Intro	duction · · · · · · · · · · · · · · · · · · ·	
7.2	2 A powerful approach to solve the trilinear optimization problem over		
	unit spheres · · · · · · · · · · · · · · · · · · ·		
7.3	Quad	Quadratic constraints · · · · · · · 154	
7.4	A spe	A special case	
7.5	Exte	nding to the biquadratic case······167	
Chapte	er 8	Conclusions	
References · · · · · · · · · · · · · · · · · · ·			
T. 1			

# Chapter 1

## Introduction

Tensor is a hot topic in the past decade. Nowadays, many real world problems can be modeled as tensor problems, just to name a few: signal processing [36,101], data analysis [17,32], chemometrics [12,13,111], hypergraph theory [34,66], diffusion magnetic resonance imaging (MRI) [3,6,44], quantum entanglement in quantum physics [35], higher order Markov chains [107] and elastic materials analysis [50,77]. Specifically, a tensor can be viewed as a multiarray: if a vector  $a = [a_1, \cdots, a_n]^T \in \mathbb{R}^n$  is a one-way-array, a matrix  $A = (a_{ij}) \in \mathbb{R}^{n \times n}$  is a two-way-array, then a tensor of order m dimension n:  $A = (a_{i_1 \cdots i_n})$  is a multiarray. Here "order" refers to the number of indices of each entry of A, e.g., a matrix is a tensor of order 2. Just imagine we have l matrices of size  $m \times n$ . By stacking them up together, we get a cube, which is a  $l \times m \times n$  tensor. If we merge k such tensors together, then we get a hyper-cube, which is a  $k \times l \times m \times n$  tensor, and so on. Particularly, if  $k = l = m = n = \cdots$ , i.e., all the dimensions of a tensor are the same, then the tensor is called a "square" tensor, which generalizes the "square" matrix and is an important type of tensors studied in this book.

We are particularly interested in the (spectral) properties of eigenvalue problems of higher order tensor and related polynomial optimization problems. In this book, the main concerns are the following five topics:

- Spectral properties and algorithms of H-eigenvalue problems;
- Spectral properties and algorithms of H-singular value problems;
- Properties and algorithms of Z-eigenvalue problems;
- Approximation methods of biquadratic optimization problems;
- Approximation methods of trilinear optimization problems.

We will introduce the related concepts and problems in the rest of this chapter.

#### 1.1 Eigenvalues problems of higher order tensors

It is well known that eigenvalues play an important role in matrix theory. For a matrix  $A \in \mathbb{C}^{n \times n}$ , if there exists a pair  $(\lambda, x) \in \mathbb{C} \times \mathbb{C}^n$  with  $x \neq 0$  such that

$$Ax = \lambda x$$

then  $\lambda$  is called an eigenvalue of A, and x is an eigenvector corresponding to  $\lambda$ . To study tensor problems, one may naturally ask a question: can we define eigenvalues and eigenvectors on tensors? The answer is positive. In 2005, the concept of eigenvalues and eigenvectors of a symmetric tensor with order even was introduced by  $Qi^{[119]}$ .

To be more specific, let  $\mathcal{A} = (a_{i_1 \cdots i_m})$  be an order m n-th dimensional real square tensor. If there is a complex number  $\lambda$  and a nonzero complex vector x that are solutions of the following homogeneous polynomial equations:

$$\mathcal{A}x^{m-1} = \lambda x^{[m-1]},\tag{1.1}$$

then  $\lambda$  is called an eigenvalue of  $\mathcal{A}$  and x the eigenvector of  $\mathcal{A}$  associated with  $\lambda$ . In problem (1.1),  $\mathcal{A}x^{m-1}$  and  $x^{[m-1]}$  are vectors, whose *i*-th entries are given by

$$(\mathcal{A}x^{m-1})_i = \sum_{i_2,\dots,i_m=1}^n a_{i,i_2\dots i_m} x_{i_2} \dots x_{i_m},$$
$$(x^{[m-1]})_i = x_i^{m-1},$$

respectively. If  $\lambda$  and x are restricted in the real field, then  $(\lambda, x)$  is called an H-eigenpaire. If an eigenvalue is not an H-eigenvalue, we call it an N-eigenvalue of  $\mathcal{A}$ . Besides the H-eigenvalues, Qi defined the Z-eigenvalues<sup>[119]</sup>: a real number  $\lambda$  and a real vector x are called Z-eigenvalue of  $\mathcal{A}$  and a Z-eigenvector of  $\mathcal{A}$  associated with the Z-eigenvalue  $\lambda$  respectively, if they are solutions of the following system:

$$\begin{cases} Ax^{m-1} = \lambda x, \\ \|x\| = 1. \end{cases} \tag{1.2}$$

Here  $\|\cdot\|$  is the Euclidean norm. If  $\lambda$  and x are complex, then they are called E-eigenvalue and E-eigenvector.

In the same year,  $\text{Lim}^{[94]}$  independently defined eigenvalues for general real tensors in the real field. In his work, the  $l^k$  eigenvalues are H-eigenvalues, while the  $l^2$  eigenvalues of tensors are the Z-eigenvalues. Note that in the case of m=2, both the H-eigenvalue and Z-eigenvalue collapse to the eigenvalues of a matrix.

The concept of singular values and singular vectors can be generalized to higher order tensors as well. Let us recall in the real case that for a matrix  $A \in \mathbb{R}^{n_1 \times n_2}$ , if there exists a triple  $(\lambda, x_1, x_2) \in \mathbb{R} \times \mathbb{R}^{n_1} \times \mathbb{R}^{n_2}$  such that

$$Ax_2 = \lambda x_1 \text{ and } A^{\mathrm{T}}x_1 = \lambda x_2,$$

then  $\lambda$  is called a singular value of A, and  $x_1$  and  $x_2$  are the left and right singular vector corresponding to  $\lambda$ , respectively.

For higher order tensors, for example, consider a third order tensor  $\mathcal{A} \in \mathbb{R}^{n_1 \times n_2 \times n_3}$ . The corresponding singular value problem can be defined as<sup>[94]</sup>

$$\begin{cases}
Ax_2x_3 = \lambda x_1, \\
Ax_1x_3 = \lambda x_2, \\
Ax_1x_2 = \lambda x_3, \\
\|x_1\| = 1, \|x_2\| = 1, \|x_3\| = 1, \lambda \in \mathbb{R}, x_1 \in \mathbb{R}^{n_1}, x_2 \in \mathbb{R}^{n_2}, x_3 \in \mathbb{R}^{n_3},
\end{cases}$$

where

$$(\mathcal{A}x_2x_3)_i = \sum_{j=1}^{n_2} \sum_{k=1}^{n_3} a_{ijk}(x_2)_j(x_3)_k,$$

$$(\mathcal{A}x_1x_3)_j = \sum_{i=1}^{n_1} \sum_{k=1}^{n_3} a_{ijk}(x_1)_i(x_3)_k,$$

$$(\mathcal{A}x_1x_2)_k = \sum_{i=1}^{n_1} \sum_{j=1}^{n_2} a_{ijk}(x_1)_i(x_2)_j.$$

For other higher order tensors, the definition is similar.

There are other kinds of eigenvalues/singular values for other types of tensors. Assume that  $p,q,n_1$  and  $n_2$  are positive integers, and  $n_1,n_2 \ge 2$ . We call  $\mathcal{A} = (a_{i_1\cdots i_p j_1\cdots j_q})$ , where  $a_{i_1\cdots i_p j_1\cdots j_q} \in \mathbb{R}$ , for  $i_k = 1, \cdots, n_1, \ k = 1, \cdots, p$ , and  $j_k = 1, \cdots, n_2, \ k = 1, \cdots, q$ , a real (p,q)-th order  $n_1 \times n_2$  dimensional rectangular tensor, or simply a real rectangular tensor. When p = q = 1,  $\mathcal{A}$  is simply a real  $n_1 \times n_2$  rectangular matrix. Denote M = p + q. If there is a complex number  $\lambda$  and two nonzero complex vectors x and y such that:

$$\left\{ \begin{array}{l} \mathcal{A}x^{p-1}y^q = \lambda x^{[M-1]}, \\ \mathcal{A}x^py^{q-1} = \lambda y^{[M-1]}, \end{array} \right.$$

where  $Ax^{p-1}y^q$  is a vector in  $\mathbb{R}^{n_1}$  whose *i*-th entry is given by

$$(\mathcal{A}x^{p-1}y^q)_i = \sum_{i_2,\dots,i_p=1}^{n_1} \sum_{j_1,\dots,j_q=1}^{n_2} a_{ii_2\dots i_p j_1\dots j_q} x_{i_2} \dots x_{i_p} y_{j_1} \dots y_{j_q},$$

and  $Ax^py^{q-1}$  is a vector in  $\mathbb{R}^{n_2}$  whose j-th entry is given by

$$(\mathcal{A}x^{p}y^{q-1})_{j} = \sum_{i_{1},\dots,i_{p}=1}^{n_{1}} \sum_{j_{2},\dots,j_{q}=1}^{n_{2}} a_{i_{1}\dots i_{p}jj_{2}\dots j_{q}} x_{i_{1}} \dots x_{i_{p}} y_{j_{2}} \dots y_{j_{q}},$$

then  $\lambda$  is called the singular value of  $\mathcal{A}$  and (x,y) are the left and right eigenvectors of  $\mathcal{A}$ , associated with  $\lambda$ , respectively. If  $(\lambda, x, y) \in \mathbb{R} \times \mathbb{R}^{n_1} \times \mathbb{R}^{n_2}$ , then  $\lambda$  is an H-singular value of  $\mathcal{A}$ , and (x,y) are the left and right H-eigenvectors associated

1.2

Th

with  $\lambda$ . If a singular value is not an H-singular value, then it is an N-singular value of  $\mathcal{A}^{[27]}$ .

Assume that  $\mathcal{A}$  is a 4-th order partially symmetric tensor, i.e., every entry of  $\mathcal{A}$  satisfies  $a_{ijkl} = a_{jikl} = a_{ijlk}$ . The M-eigenvalue system of  $\mathcal{A}$  is defined as

$$\begin{cases} Axyy = \lambda x, \\ Axxy = \lambda y, \\ \|x\| = 1, \|y\| = 1, x \in \mathbb{R}^{n_1}, y \in \mathbb{R}^{n_2}. \end{cases}$$

Here  $\lambda \in \mathbb{R}$  is called the M-eigenvalue of  $\mathcal{A}$ . This type of eigenvalue was introduced by Qi et al. <sup>[122]</sup>, which finds applications in nonlinear elastic materials analysis <sup>[122]</sup> and quantum entanglement problems in quantum physics <sup>[35]</sup>.

There are other type of eigenvalues, e.g, the  $Z_1$ -eigenvalue of a transition probability tensor was introduced by Li and and  $Ng^{[90]}$ , where a transition probability tensor  $\mathcal{P}$  of order m dimension n is defined to satisfy  $\sum_{i_1=1}^n p_{i_1 i_2 \cdots i_m} = 1$ , and the  $Z_1$ -eigenvalue problem is replacing the 2-norm in the Z-eigenvalue problem (1.2) by the 1-norm, i.e., the  $Z_1$ -eigenvalue problem is defined by

$$\begin{cases} \mathcal{A}x^{m-1} = \lambda x, \\ \|x\|_1 = 1. \end{cases} \tag{1.3}$$

This type of eigenvalue problems can be used to approximate the limit probability distribution vector of a higher order Markov chain<sup>[90]</sup>.

#### 1.2 Related polynomial optimization problems

Before proceeding, we first introduce the symmetric property of higher order tensors. A tensor  $\mathcal{A}$  is called symmetric, if every of its entry  $a_{i_1 \cdots i_m}$  is invariant under all permutations of  $\{i_1, \cdots, i_m\}$ , i.e., the following equalities always hold:

$$a_{i_1\cdots i_m}=a_{\sigma(i_1\cdots i_m)},$$

where  $\sigma(\cdot)$  is the permutation function.

Every symmetric tensor is related to a polynomial function uniquely<sup>[33]</sup>. Specifically, suppose that  $\mathcal{A}$  is an order m n-th dimensional symmetric tensor. Then

$$f_{\mathcal{A}}(x) = \mathcal{A}x^m = \sum_{i_1, \dots, i_m = 1}^n a_{i_1 \dots i_m} x_{i_1} \dots x_{i_m}$$

is the associated polynomial function.

Consider the following homogenous polynomial optimization problem of degree m:

$$\max \quad \mathcal{A}x^m,$$
 s.t. 
$$||x|| = 1, x \in \mathbb{R}^n.$$
 (1.4)

The corresponding Lagrangian function of (1.4) is given by

$$L(x,\lambda) = \mathcal{A}x^m - \lambda(x^{\mathrm{T}}x - 1),$$

where  $\lambda \in \mathbb{R}$  is the Lagrangian multiplier. Taking the gradient of  $L(\cdot, \cdot)$  with respect to x and  $\lambda$  respectively and setting them both to zero, we get the following KKT system of (1.4)

$$\begin{cases} Ax^{m-1} = \frac{2\lambda}{m}x, \\ x^{\mathrm{T}}x = 1, \end{cases}$$

which is in fact the Z-eigenvalue problem of  $\mathcal{A}$  if we let  $\lambda := 2\lambda/m$  in the above system. Thus finding the largest Z-eigenvalue of  $\mathcal{A}$  is equivalent to solving the polynomial optimization problem (1.4).

Suppose the symmetric tensor  $\mathcal{A}$  is of even order, order Then the KKT system of the following problem

$$\max \quad \mathcal{A}x^m,$$
  
s.t. 
$$||x||_m = 1, x \in \mathbb{R}^n$$

is

$$\begin{cases} Ax^{m-1} = \lambda x^{[m-1]}, \\ \sum_{i=1}^{n} x_i^m = 1, \end{cases}$$

where  $||x||_m = (\sum_{i=1}^n |x_i|^m)^{\frac{1}{m}}$ . Note that the first formula is homogenous with respect to x. Thus the second formula can be eliminated, which is the H-eigenvalue problem of  $\mathcal{A}$ .

Finding the largest M-eigenvalue of a 4-th order partially symmetric tensor  $\mathcal{A}$  is equivalent to solving the following homogeneous polynomial problem with degree 4:

$$\max \quad \mathcal{A}xxyy = \sum_{i,j=1}^{n_1} \sum_{k,l=1}^{n_2} a_{ijkl} x_i x_j y_k y_l,$$
s.t. 
$$||x|| = 1, ||y|| = 1, x \in \mathbb{R}^{n_1}, y \in \mathbb{R}^{n_2}.$$
(1.5)

This is known as the biquadratic optimization problem  $^{[62,97,122]}$ .

Similarly, finding the largest singular value of a tensor of order d is equivalent to solving the following multi-linear optimization problem:

$$\max \quad \mathcal{A}x_1 \cdots x_d = \sum_{i_1=1}^{n_1} \cdots \sum_{i_d=1}^{n_d} a_{i_1 \cdots i_d}(x_1)_{i_1} \cdots (x_d)_{i_d},$$
s.t. 
$$\|x^1\| = 1, \cdots, \|x^d\| = 1, x^i \in \mathbb{R}^{n_i}, i = 1, \cdots, d.$$
(1.6)

The polynomial optimization in binary variables or in mixed variables was studied in [54,56]:

$$\max_{\mathbf{x}} \quad \mathcal{A}x^m, 
\text{s.t.} \quad x \in \{-1, 1\}^n,$$
(1.7)

$$\max \quad \mathcal{A}x_1 \cdots x_d, \text{s.t.} \quad x_i \in \{-1, 1\}^{n_i}, i = 1, \cdots, d,$$
 (1.8)

and

max 
$$\mathcal{A}x_1 \cdots x_d$$
,  
s.t.  $||x_i|| = 1, x_i \in \mathbb{R}^{n_i}, i = 1, \cdots, l$ , (1.9)  
 $x_j \in \{-1, 1\}^{n_j}, j = l + 1, \cdots, d$ ,

which relate to the  $\infty \to 1$ -norm of higher order tensors<sup>[56]</sup>.

The biquadratic optimization over quadratic constraints or standard simplex was studied by Zhang et al. [169], Ling et al. [98] and Bomze et al. [7]:

max 
$$\mathcal{A}xxyy$$
,  
s.t.  $x^{\mathrm{T}}A^{i}x \leq 1$ ,  $i = 1, \dots, p$ ,  $y^{\mathrm{T}}B^{j}y \leq 1$ ,  $j = 1, \dots, q$ , 
$$(1.10)$$

max 
$$\mathcal{A}xxyy$$
,  
s.t. 
$$\sum_{i=1}^{m} x_i = 1, x \ge 0,$$

$$\sum_{j=1}^{n} y_j = 1, y \ge 0,$$
(1.11)

which arise from portfolio selection<sup>[98]</sup>.

#### 1.3 Applications

This section lists some applications of eigenvalues and singular values of tensors and their related polynomial optimization problems.

Application of the H-eigenvalue to maximum clique problem $^{[14,15]}$ :

A k-uniform hypergraph, or simply a k-graph, is a pair G=(V,E), where  $V=\{1,\cdots,n\}$  is a finite set of vertices and  $E\subseteq \binom{V}{k}$  is a set of k-subsets of V, each of which is called a hyperedge. Particularly, a 2-graph is the usual graph. A subset

1.3 Applications 7

of vertices  $C \subseteq V$  is called a hyperclique if  $\binom{C}{k} \subseteq E$ . A clique is said to be maximal if it is not contained in any other clique, while it is called maximum if it has maximum cardinality. Specifically, the maximal clique problem of a 2-graph is to find the largest subset of vertices in which each point is directly connected to every other vertex in the subset. The clique number of a graph G, denoted by G, is defined as the cardinality of a maximum clique. Computing G0 is NP-hard. Bulò and Pelillo<sup>[15]</sup> provided a new upper bound for G0 of an undirected graph G3.

$$\omega(G) \leqslant \frac{\rho(\mathcal{A})}{k!} + k,$$

where  $\mathcal{H}$  is the induced tensor of a k-clique (k+1)-graph constructed from G, and  $\rho(\mathcal{H})$  is the largest H-eigenvalue of  $\mathcal{H}$ . A new lower bound was also provided via the largest H-eigenvalue.

Application of singular value and Z-eigenvalue to the best rank-one approximation of higher order tensors<sup>[36,78]</sup>.

In signal processing, data analysis, chemometrics, and psychology, data may sometimes be organized as higher order higher dimension tensors. How to extract useful information from this data is an important problem<sup>[36]</sup>. One possible way is to find a rank-one tensor which is the closest one to the original tensor under the Euclidean distance, which contains most of the information of the original tensor. To be more specific, let the original d-th order tensor be  $\mathcal{F}$  and the rank-one tensor be  $\mathcal{B} = x^1 \otimes x^2 \otimes \cdots \otimes x^d$ , where " $\otimes$ " denotes the Kronecker product, and  $b_{i_1 \cdots i_d} = x_{i_1}^1 \cdots x_{i_d}^d$ . The associated problem can be written as:

min 
$$\|\mathcal{F} - \mathcal{B}\|_F = \|\mathcal{F} - x_1 \otimes \cdots \otimes x_d\|_F$$
,  
s.t.  $x_i \in \mathbb{R}^{n_i}, i = 1, \cdots, d$ ,

where  $\|\mathcal{F}\|_F = \sqrt{\sum_{i_1=1}^{n_1} \cdots \sum_{i_d=1}^{n_d} f_{i_1\cdots i_d}^2}$ . If we restrict the norm of  $x^i$  to be one, then it can be rewritten as

min 
$$\|\mathcal{F} - \lambda x_1 \otimes \cdots \otimes x_d\|_F$$
,  
s.t.  $\lambda \in \mathbb{R}, \|x_i\| = 1, x_i \in \mathbb{R}^{n_i}, i = 1, \cdots, d$ .

Note that

$$\|\mathcal{F} - \lambda x_1 \otimes \cdots \otimes x_d\|_F^2$$

$$= \|\mathcal{F}\|_F^2 - 2\lambda \cdot \mathcal{F} x_1 \cdots x_d + \lambda^2 \|x_1 \otimes \cdots \otimes x_d\|_F^2$$

$$= \lambda^2 - 2\lambda \cdot \mathcal{F} x_1 \cdots x_d + \|\mathcal{F}\|_F^2.$$

Minimizing this with respect to  $\lambda$ , we get

$$\|\mathcal{F} - \lambda x_1 \otimes \cdots \otimes x_d\|_F^2 = \|\mathcal{F}\|_F^2 - (\mathcal{F} x_1 \cdots x_d)^2,$$

which reduces to

$$\max \quad |\mathcal{F}x_1 \cdots x_d| = \mathcal{F}x_1 \cdots x_d,$$
s.t. 
$$||x_1|| = 1, \cdots, ||x_d|| = 1, x^i \in \mathbb{R}^n, i = 1, \cdots, d,$$

which reduces to finding the largest singular value of  $\mathcal{F}$ .

Suppose the symmetric tensor  $\mathcal{A}$  is of order m dimension n. The best symmetric rank-one approximation to  $\mathcal{A}$  is

min 
$$\|\mathcal{A} - \lambda x \otimes \cdots \otimes x\|_F$$
,  
s.t.  $\lambda \in \mathbb{R}, \|x\| = 1, x \in \mathbb{R}^n$ .

Using a similar argument, the problem is equivalent to

$$\max \quad |\mathcal{A}x^m|,$$
  
s.t. 
$$||x|| = 1,$$

which is finding the extreme Z-eigenvalue of A.

In automatic control, one needs to determine whether a quartic form is positive semidefinite or not<sup>[110]</sup>. Specifically, for any  $x \in \mathbb{R}^n$ , to determine whether

$$\sum_{i,j,k,l=1}^{n} a_{ijkl} x_i x_j x_k x_l \geqslant 0$$

or not. by restricting x on the unit sphere, it can be written as the following problem:

min 
$$\mathcal{A}x^4 := \sum_{i,j,k,l=1}^n a_{ijkl}x_ix_jx_kx_l,$$
  
s.t.  $||x|| = 1, x \in \mathbb{R}^n,$ 

which is finding the smallest Z-eigenvalue of A.

The M-eigenvalue finds applications in nonlinear elastic materials analysis and quantum physics. In nonlinear elastic materials where one needs to determines whether the strong or ordinary ellipticity condition holds or not, where the strong one holds if and only if the minimum of a biquadratic function over unit spheres in  $\mathbb{R}^3$  is positive:

$$\min \left\{ f(x,y) = \sum_{i,j,k,l=1}^{3} a_{ijkl} x_i x_j y_k y_l \, \middle| \, ||x|| = 1, ||y|| = 1 \right\} > 0,$$

and the ordinary one holds if and only if the minimum is nonnegative. In quantum physics, where a quantum state is separable or entangled is a fundamental and